

Categorification of Algebras: 2-Algebras

İbrahim İlker Akça¹ , Ummahan Ege Arslan² 

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2-Categories,
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Abstract – This paper introduces a categorification of k -algebras called 2-algebras, where k is a commutative ring. We define the 2-algebras as a 2-category with single object in which collections of all 1-morphisms and all 2-morphisms are k -algebras. It is shown that the category of 2-algebras is equivalent to the category of crossed modules in commutative k -algebras.

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1. Introduction

The term “categorification” coined by Louis Crane refers to the process of replacing set theoretic concepts by category-theoretic analogues in mathematics. A categorified version of a group is a 2-group. Internal categories in the category of groups are exactly the same as 2-groups. The Brown-Spencer theorem [3] thus constructs the associated 2-group of a crossed module given by Whitehead [11] to define an algebraic model for a “(connected) homotopy 2-type”. The fact that the composition in the internal category must be a group homomorphism implies that the “interchange law” must hold. This equation is in fact equivalent via the Brown-Spencer result to the Peiffer identity.

We will concern in this paper exclusively with categorification of algebras. We will obtain analogous results in (commutative) algebras with regard to Porter’s work [9]. He states that there is an equivalence of categories between the category of internal categories in the category of k -algebras and the category of crossed modules of commutative k -algebras. Since the internal category in the category of k -algebras is a categorification of k -algebras, this internal category will be called as “strict 2-algebra” in this work. We define the strict 2-algebra by means of 2-module being a category in the category of modules as a 2-category with single object in which collections of 1-morphisms and 2-morphisms are k -algebras and we denote the category of strict 2-algebras by **2Alg**. Given a group G , it is known that automorphisms of G yield a 2-group. Analogous result in commutative algebras can be given that multiplications of C yield a strict 2-algebra where C is a commutative R -algebra and R is a commutative k -algebra.

A crossed module $\mathcal{A} = (\partial : C \longrightarrow R)$ of commutative algebras is given by an algebra morphism $\partial : C \longrightarrow R$

¹iakca@ogu.edu.tr (Corresponding Author); ²uege@ogu.edu.tr

^{1,2}Department of Mathematics and Computer Sciences, Faculty of Science, Eskişehir Osmangazi University, Eskişehir, Turkey

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together with an action \cdot of R on C such that the relations below hold for each $r \in R$ and each $c, c' \in C$,

$$\begin{aligned} \partial(r \cdot c) &= r\partial(c) \\ \partial(c) \cdot c' &= cc'. \end{aligned}$$

In this paper we show that the category of strict 2-algebras is equivalent to the category of crossed modules in commutative algebras.

2. Internal Categories and 2-categories

We begin by recalling internal categories as well as 2-categories. Ehresmann defined internal categories in [5], and by now they are an important part of category theory [4].

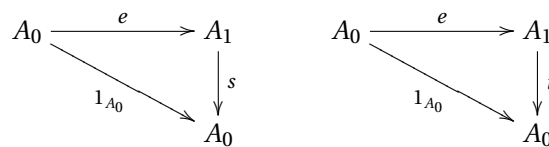
2.1. Internal categories

Definition 2.1. Let \mathbf{C} be any category. An internal category in \mathbf{C} , say \mathbf{A} , consists of:

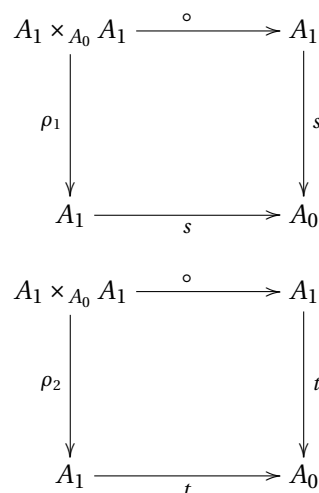
- an object of objects $A_0 \in \mathbf{C}$
- an object of morphisms $A_1 \in \mathbf{C}$,

together with

- source and target morphisms $s, t : A_1 \rightarrow A_0$,
- an identity-assigning morphism $e : A_0 \rightarrow A_1$,
- a composition morphism $\circ : A_1 \times_{A_0} A_1 \rightarrow A_1$ such that the following diagrams commute, expressing the usual category laws:
- laws specifying the source and target of identity morphisms:



- laws specifying the source and target of composite morphisms:



- the associative law for composition of morphisms:

$$\begin{array}{ccc}
 A_1 \times_{A_0} A_1 \times_{A_0} A_1 & \xrightarrow{\circ} & A_1 \times_{A_0} A_1 \\
 \downarrow \rho_2 & & \downarrow t \\
 A_1 \times_{A_0} A_1 & \xrightarrow{t} & A_0
 \end{array}$$

- the left and right unit laws for composition of morphisms:

$$\begin{array}{ccccc}
 A_0 \times_{A_0} A_1 & \xrightarrow{e \times_{A_0} 1_{A_1}} & A_1 \times_{A_0} A_1 & \xleftarrow{1_{A_1} \times_{A_0} e} & A_1 \times_{A_0} A_0 \\
 & \searrow \rho_2 & \downarrow \circ & \swarrow \rho_1 & \\
 & & A_1 & &
 \end{array}$$

Here, the pullback $A_1 \times_{A_0} A_1$ is defined via the square:

$$\begin{array}{ccc}
 A_1 \times_{A_0} A_1 & \xrightarrow{\rho_2} & A_1 \\
 \rho_1 \downarrow & & \downarrow s \\
 A_1 & \xrightarrow{t} & A_0.
 \end{array}$$

We denote this internal category with $A = (A_0, A_1, s, t, e, \circ)$.

Definition 2.2. Let \mathbf{C} be a category. Given internal categories A and A' in \mathbf{C} , an **internal functor** between them, say $F : A \rightarrow A'$, consists of

- a morphism $F_0 : A_0 \rightarrow A'_0$,
- a morphism $F_1 : A_1 \rightarrow A'_1$

such that the following diagrams commute, corresponding to the usual laws satisfied by a functor:

- preservation of source and target:

$$\begin{array}{ccc}
 A_1 \xrightarrow{s} A_0 & & A_1 \xrightarrow{t} A_0 \\
 F_1 \downarrow & & F_1 \downarrow \\
 A'_1 \xrightarrow{s'} A'_0 & & A'_1 \xrightarrow{t'} A'_0 \\
 \downarrow F_0 & & \downarrow F_0
 \end{array}$$

- preservation of identity morphisms:

$$\begin{array}{ccc}
 A_0 \xrightarrow{e} A_1 & & \\
 F_0 \downarrow & & \downarrow F_1 \\
 A'_0 \xrightarrow{e'} A'_1 & &
 \end{array}$$

- preservation of composite morphisms:

$$\begin{array}{ccc}
 A_1 \times_{A_0} A_1 & \xrightarrow{F_1 \times_{A_0} F_1} & A'_1 \times_{A'_0} A'_1 \\
 \downarrow \circ & & \downarrow \circ' \\
 A_1 & \xrightarrow{F_1} & A'_1
 \end{array}$$

Given two internal functors $F : A \rightarrow A'$ and $G : A' \rightarrow A''$ in some category \mathbf{C} , we define their composite $FG : A \rightarrow A''$ by taking $(FG)_0 = F_0 G_0$ and $(FG)_1 = F_1 G_1$. Similarly, we define the identity internal functor in \mathbf{C} , $1_A : A \rightarrow A$ by taking $(1_A)_0 = 1_{A_0}$ and $(1_A)_1 = 1_{A_1}$.

Definition 2.3. Let \mathbf{C} be a category. Given two internal functors $F, G : A \rightarrow A'$ in \mathbf{C} , an **internal natural transformation** in \mathbf{C} between them, say $\theta : F \Rightarrow G$, is a morphism $\theta : A_0 \rightarrow A'_1$ for which the following diagrams commute, expressing the usual laws satisfied by a natural transformation:

- laws specifying the source and target of a natural transformation:

$$\begin{array}{ccc}
 A_0 & \xrightarrow{\theta} & A'_1 \\
 & \searrow F_0 & \downarrow s' \\
 & & A'_0
 \end{array}
 \qquad
 \begin{array}{ccc}
 A'_0 & \xrightarrow{\theta} & A'_1 \\
 & \searrow G_0 & \downarrow t' \\
 & & A_0
 \end{array}$$

- the commutative square law:

$$\begin{array}{ccc}
 A_1 & \xrightarrow{\Delta(s\theta \times G)} & A'_1 \times_{A'_0} A'_1 \\
 \downarrow \Delta(F \times t\theta) & & \downarrow \circ' \\
 A'_1 \times_{A'_0} A'_1 & \xrightarrow{\circ'} & A'_1
 \end{array}$$

Given an internal functor $F : A \rightarrow A'$ in \mathbf{C} , the identity internal natural transformation $1_F : F \Rightarrow F$ in \mathbf{C} is given by $1_F = F_0 e$.

2.2. 2-categories

Definition 2.4. A 2-category \mathcal{G} consists of a class of objects G_0 and for any pair of objects (A, B) a small category of morphisms $\mathcal{G}(A, B)$ -with objects $G_1(A, B)$ and morphisms $G_2(A, B)$ -, along with composition functors

$$\bullet : \mathcal{G}(A, B) \times \mathcal{G}(B, C) \rightarrow \mathcal{G}(A, C)$$

for every triple (A, B, C) of objects and identity functors from the terminal category to $\mathcal{G}(A, A)$

$$i_A : 1 \rightarrow \mathcal{G}(A, A)$$

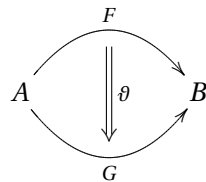
for all objects A such that \bullet is associative and

$$F \bullet i_B = F = i_A \bullet F \quad \text{as well as} \quad \vartheta \bullet I_{i_B} = \vartheta = I_{i_A} \bullet \vartheta$$

hold for all $F \in G_1(A, B)$ and $\vartheta \in G_2(A, B)$ where source and target morphisms are defined by

$$\begin{array}{ccc}
 A & \xrightarrow{F} & B \\
 \\
 s: G_1(A, B) & \longrightarrow & G_0 \\
 F & \longmapsto & s(F) = A \\
 \\
 t: G_1(A, B) & \longrightarrow & G_0 \\
 F & \longmapsto & t(F) = B
 \end{array}$$

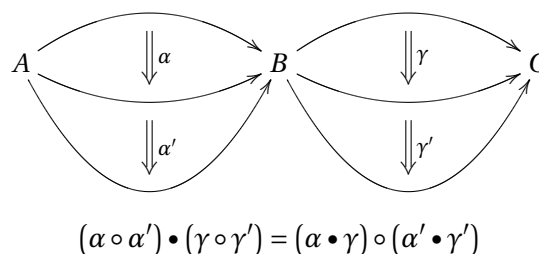
for $F \in G_1(A, B)$ and



$$\begin{array}{ccc}
 s: G_2(A, B) & \longrightarrow & G_1 \\
 \vartheta & \longmapsto & s(\vartheta) = F \\
 \\
 t: G_2(A, B) & \longrightarrow & G_0 \\
 \vartheta & \longmapsto & t(\vartheta) = G
 \end{array}$$

for $\vartheta : F \longrightarrow G \in G_2(A, B)$. For all pairs of objects (A, B) elements of $G_1(A, B)$ are called 1-morphisms or 1-cells of \mathcal{G} and elements of $G_2(A, B)$ are called 2-morphisms or 2-cells of \mathcal{G} . We write G_1 and G_2 for the classes of all 1-morphisms and 2-morphisms respectively.

There are two ways of composing 2-morphisms: using the composition \circ inside the categories $\mathcal{G}(A, B)$, called vertical composition, and using the morphism level of the functor \bullet , called horizontal composition. These compositions must satisfy the following equation: for $\alpha, \alpha' \in G_2(A, B)$ with $t(\alpha) = s(\alpha')$ and $\gamma, \gamma' \in G_2(B, C)$ with $t(\gamma) = s(\gamma')$



which is called “interchange law”.

3. Constructions of Two-Algebras

In this section we will construct 2-algebras by categorification. We can categorify the notion of an algebra by replacing the equational laws by isomorphisms satisfying extra structure and properties we expect. In [2]

Baez and Crans introduce the Lie 2-algebra by means of the concept of 2-vector space defined as an internal category in the category of vector spaces by them. Obviously we get a new notion of “2-module” which can be considered as an internal category in the category of modules and we categorify the notion of an algebra.

3.1. 2-Modules

A categorified module or “2-module” should be a category with structure analogous to that of a k -module, with functors replacing the usual k -module operations. Here we instead define a 2-module to be an internal category in a category of k -modules \mathbf{Mod} . Since the main component part of a k -algebra is a k -module, a 2-algebra will have an underlying 2-module of this sort. In this section we thus first define a category of these 2-modules.

In the rest of this paper, the terms a module and an algebra will always refer to a k -module and a k -algebra.

Definition 3.1. A 2-module is an internal category in \mathbf{Mod} .

Thus, a 2-module M is a category with a module of objects M_0 and a module of morphisms M_1 , such that the source and target maps $s, t : M_1 \rightarrow M_0$, the identity assigning map $e : M_0 \rightarrow M_1$, and the composition map $\circ : M_1 \times_{M_0} M_1 \rightarrow M_1$ are all module morphisms. We write a morphism as $a : x \rightarrow y$ when $s(a) = x$ and $t(a) = y$, and sometimes we write $e(x)$ as 1_x .

The following proposition is given for the \mathbf{Vect} vector space category in [2]. But we rewrite this proposition for \mathbf{Mod} .

Proposition 3.2. It is defined a 2-module by specifying the modules M_0 and M_1 along with the source, target and identity module morphisms and the composition morphism \circ , satisfying the conditions of Definition 2.1. The composition map is uniquely determined by

$$\begin{aligned} \circ : M_1 \times_{M_0} M_1 &\longrightarrow M_1 \\ (a, b) &\longmapsto \circ(a, b) = a \circ b = a + b - (es)(b). \end{aligned}$$

Proof.

First given modules M_0, M_1 and module morphisms $s, t : M_1 \rightarrow M_0$ and $e : M_0 \rightarrow M_1$, we will define a composition operation that satisfies the laws in the definition of internal category, obtaining a 2-module.

Given $a, b \in M_1$ such that $t(a) = s(b)$, i.e.

$$a : x \rightarrow y \text{ and } b : y \rightarrow z$$

we define their composite \circ by

$$\begin{aligned} \circ : M_1 \times_{M_0} M_1 &\longrightarrow M_1 \\ (a, b) &\longmapsto \circ(a, b) = a \circ b = a + b - (es)(b). \end{aligned}$$

We will show that with this composition \circ the diagrams of the definition of internal category commute. The triangles specifying the source and target of the identity-assigning morphism do not involve composition.

The second pair of diagrams commute since

$$\begin{aligned}
 s(a \circ b) &= s(a + b - (es)(b)) \\
 &= s(a) + s(b) - (se)(s(b)) \\
 &= s(a) + s(b) - s(b) \\
 &= s(a) = x
 \end{aligned}$$

and since $t(a) = s(b)$,

$$\begin{aligned}
 t(a \circ b) &= t(a + b - (es)(b)) \\
 &= t(a) + t(b) - (te)(s(b)) \\
 &= t(a) + t(b) - s(b) \\
 &= t(b) = z.
 \end{aligned}$$

The associative law holds for composition because module addition is associative. Finally the left and right unit laws are satisfied since given $a : x \rightarrow y$,

$$\begin{aligned}
 e(x) \circ a &= e(x) + a - (es)(a) \\
 &= e(x) + a - e(x) \\
 &= a
 \end{aligned}$$

and

$$\begin{aligned}
 a \circ e(y) &= a + e(y) - (es)(e(y)) \\
 &= a + e(y) - e(y) \\
 &= a.
 \end{aligned}$$

We thus have a 2-module.

Given a 2-module M , we shall show that its composition must be defined by the formula given above. Suppose that (a, g) and (a', g') are composable pairs of morphisms in M_1 , i.e.

$$a : x \rightarrow y \text{ and } b : y \rightarrow z$$

and

$$a' : x' \rightarrow y' \text{ and } b' : y' \rightarrow z'.$$

Since the source and target maps are module morphisms, $(a + a', b + b')$ also forms a composable pair, and since that the composition is module morphism

$$(a + a') \circ (b + b') = a \circ b + a' \circ b'.$$

Then if (a, b) is a composable pair, i.e, $t(a) = s(b)$, we have

$$\begin{aligned}
 a \circ b &= (a + 1_{M_1}) \circ (1_{M_1} + b) \\
 &= (a + e(s(b) - s(b))) \circ (e(s(b) - s(b)) + b) \\
 &= (a - e(s(b)) + e(s(b))) \circ (e(s(b)) - e(s(b)) + b) \\
 &= (a \circ e(s(b))) + (-e(s(b)) + e(s(b))) \circ (-e(s(b)) + b) \\
 &= a \circ e(s(b)) + (-e(s(b)) \circ (-e(s(b)))) + (e(s(b)) \circ b) \\
 &= a - e(s(b)) + b \\
 &= a + b - e(s(b)).
 \end{aligned}$$

This show that we can define \circ by

$$\begin{aligned}
 \circ : M_1 \times_{M_0} M_1 &\longrightarrow M_1 \\
 (a, b) &\longmapsto \circ(a, b) = a \circ b = a + b - e(s(b)).
 \end{aligned}$$

Corollary 3.3. For $b \in \ker s$, we have

$$\begin{aligned}
 a \circ b &= a + b - (es)(b) \\
 &= a + b.
 \end{aligned}$$

Definition 3.4. Let M and N be 2-modules, a 2-module functor $F : M \longrightarrow N$ is an internal functor in **Mod** from M to N . 2-modules and 2-module functors between them is called the category of 2-modules denoted by **2Mod**.

After we get the definition of a 2-module, we define the definition of a categorified algebra which is main concept of this paper.

3.2. Two-algebras

Definition 3.5. A weak 2-algebra consists of

- a 2-module A equipped with a functor $\bullet : A \times A \longrightarrow A$, which is defined by $(x, y) \mapsto x \bullet y$ and bilinear on objects and defined by $(f, g) \mapsto f \bullet g$ on morphisms satisfying interchange law, i.e.,

$$(f_1 \bullet g_1) \circ (f_2 \bullet g_2) = (f_1 \circ f_2) \bullet (g_1 \circ g_2)$$

- k -bilinear natural isomorphisms

$$\alpha_{x,y,z} : (x \bullet y) \bullet z \longrightarrow x \bullet (y \bullet z)$$

$$l_x : 1 \bullet x \longrightarrow x$$

$$r_x : x \bullet 1 \longrightarrow x$$

such that the following diagrams commute for all objects $w, x, y, z \in A_0$.

$$\begin{array}{ccc}
 ((w \bullet x) \bullet y) \bullet z & \xrightarrow{\alpha_{w \bullet x, y, z}} & (w \bullet x) \bullet (y \bullet z) \\
 \alpha_{w, x, y} \bullet 1_z \downarrow & & \searrow \alpha_{w, x, y \bullet z} \\
 (w \bullet (x \bullet y)) \bullet z & \xrightarrow{\alpha_{w, x \bullet y, z}} & w \bullet ((x \bullet y) \bullet z) \xrightarrow{1_w \bullet \alpha_{x, y, z}} w \bullet (x \bullet (y \bullet z))
 \end{array}$$

$$\begin{array}{ccc}
 (x \bullet 1) \bullet y & \xrightarrow{\alpha_{x,1,y}} & x \bullet (1 \bullet y) \\
 & \searrow r_x \bullet 1_y & \downarrow 1_x \bullet l_y \\
 & & x \bullet y
 \end{array}$$

A strict 2-algebra is the special case where $\alpha_{x,y,z}$, l_x , r_x are all identity morphisms. In this case we have

$$(x \bullet y) \bullet z = x \bullet (y \bullet z)$$

$$1 \bullet x = x, x \bullet 1 = x$$

Strict 2-algebra is called commutative strict 2-algebra if $x \bullet y = y \bullet x$ for all objects $x, y \in A_0$ and $f \bullet g = g \bullet f$ for all morphisms $f, g \in A_1$.

In the rest of this paper, the term 2-algebra will always refer to a commutative strict 2-algebra. A homomorphism between 2-algebras should preserve both the 2-module structure and the \bullet functor.

Definition 3.6. Given 2-algebras A and A' , a homomorphism

$$F: A \longrightarrow A'$$

consists of

- a linear functor F from the underlying 2-module of A to that of A' , and
- a bilinear natural transformation

$$F_2(x, y) : F_0(x) \bullet F_0(y) \longrightarrow F_0(x \bullet y)$$

- an isomorphism $F : 1' \longrightarrow F_0(1)$ where 1 is the identity object of A and $1'$ is the identity object of A' , such that the following diagrams commute for $x, y, z \in A_0$,

$$\begin{array}{ccccc}
 (F(x) \bullet F(y)) \bullet F(z) & \xrightarrow{F_2 \bullet 1} & F(x \bullet y) \bullet F(z) & \xrightarrow{F_2} & F((x \bullet y) \bullet z) \\
 \alpha_{F(x), F(y), F(z)} \downarrow & & & & \downarrow F(\alpha_{x,y,z}) \\
 F(x) \bullet (F(y) \bullet F(z)) & \xrightarrow{1 \bullet F_2} & F(x) \bullet F(y \bullet z) & \xrightarrow{F_2} & F(x \bullet (y \bullet z)).
 \end{array}$$

$$\begin{array}{ccc}
 1' \bullet F(x) & \xrightarrow{l'_{F(x)}} & F(x) \\
 F_0 \bullet 1 \downarrow & & \uparrow F(l_x) \\
 F(1) \bullet F(x) & \xrightarrow{F_2} & F(1 \bullet x).
 \end{array}$$

$$\begin{array}{ccc}
 F(x) \bullet 1' & \xrightarrow{r'_{F(x)}} & F(x) \\
 1 \bullet F_0 \downarrow & & \uparrow F(r_x) \\
 F(x) \bullet F(1) & \xrightarrow{F_2} & F(x \bullet 1).
 \end{array}$$

Definition 3.7. 2-algebras and homomorphisms between them give the category of 2-algebras denoted by

2Alg.

Therefore if $A = (A_0, A_1, s, t, e, \circ, \bullet)$ is a 2-algebra, A_0 and A_1 are algebras with this \bullet bilinear functor. Thus we can take that 2-algebra is a 2-category with a single object say $*$, and A_0 collections of its 1-morphisms and A_1 collections of its 2-morphisms are algebras with identity.

3.3. Multiplication Algebras yield a 2-algebra

In [8] Norrie developed Lue's work [6] and introduced the notion of an actor of crossed modules of groups where it is shown to be the analogue of the automorphism group of a group. In the category of commutative algebras the appropriate replacement for automorphism groups is the multiplication algebra $\mathcal{M}(C)$ of an algebra C which is defined by MacLane [7].

Let C be an associative (not necessarily unitary or commutative) R -algebra. We recall Mac Lane's construction of the R -algebra $\text{Bim}(C)$ of bimultipliers of C [7].

An element of $\text{Bim}(C)$ is a pair (γ, δ) of R -linear mappings from C to C such that

$$\gamma(cc') = \gamma(c)c'$$

$$\delta(cc') = c\delta(c')$$

and

$$c\gamma(c') = \delta(c)c'.$$

$\text{Bim}(C)$ has an obvious R -module structure and a product

$$(\gamma, \delta)(\gamma', \delta') = (\gamma\gamma', \delta'\delta),$$

the value of which is still in $\text{Bim}(C)$.

Suppose that $\text{Ann}(C) = 0$ or $C^2 = C$. Then $\text{Bim}(C)$ acts on C by

$$\begin{aligned} \text{Bim}(C) \times C &\rightarrow C; & ((\gamma, \delta), c) &\mapsto \gamma(c), \\ C \times \text{Bim}(C) &\rightarrow C; & (c, (\gamma, \delta)) &\mapsto \delta(c) \end{aligned}$$

and there is a

$$\begin{aligned} \mu: C &\longrightarrow \text{Bim}(C) \\ c &\longmapsto (\gamma_c, \delta_c) \end{aligned}$$

with

$$\gamma_c(x) = cx \quad \text{and} \quad \delta_c(x) = xc.$$

Commutative case: we still assume $\text{Ann}(C) = 0$ or $C^2 = C$. If C is a commutative R -algebra and $(\gamma, \delta) \in \text{Bim}(C)$, then $\gamma = \delta$. This is because for every $x \in C$:

$$\begin{aligned} x\delta(c) &= \delta(c)x = c\gamma(x) = \gamma(x)c \\ &= \gamma(xc) = \gamma(cx) = \gamma(c)x = x\gamma(c). \end{aligned}$$

Thus $\text{Bim}(C)$ may be identified with the R -algebra $\mathcal{M}(C)$ of multipliers of C . Recall that a multiplier of C is

a linear mapping $\lambda : C \rightarrow C$ such that for all $c, c' \in C$

$$\lambda(cc') = \lambda(c)c'.$$

Also $\mathcal{M}(C)$ is commutative as

$$\lambda'\lambda(xc) = \lambda'(\lambda(x)c) = \lambda(x)\lambda'(c) = \lambda'(c)\lambda(x) = \lambda\lambda'(cx) = \lambda\lambda'(xc)$$

for any $x \in C$. Thus $\mathcal{M}(C)$ is the set of all multipliers λ such that $\lambda\gamma = \gamma\lambda$ for every multiplier γ .

In [10] Porter states that automorphisms of a group G yield a 2-group. The appropriate analogue of this result in algebra case can be given. We claim that multiplications of an R -algebra C give a 2-algebra which is called a multiplication 2-algebra.

Let k be a commutative ring, R be a k -algebra with identity and C be a commutative R -algebra with $Ann(C) = 0$ or $C^2 = C$. Take $A_0 = \mathcal{M}(C)$ and say 1-morphisms to the elements of A_0 . We define the action of $\mathcal{M}(C)$ on C as follows:

$$\begin{aligned} \mathcal{M}(C) \times C &\longrightarrow C \\ (f, x) &\longmapsto f \blacktriangleright x = f(x). \end{aligned}$$

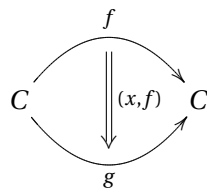
Using the action of $\mathcal{M}(C)$ on C , we can form the semidirect product

$$C \rtimes \mathcal{M}(C) = \{(x, f) | x \in C, f \in \mathcal{M}(C)\}$$

with multiplication

$$(x, f)(x', f') = (f \blacktriangleright x' + f' \blacktriangleright x + x'x, f'f).$$

Take $A_1 = C \rtimes \mathcal{M}(C)$ and say 2-morphisms to the elements of A_1 . Therefore we get the following diagram for $(x, f) \in C \rtimes \mathcal{M}(C)$,



and we define the source, target and identity assigning maps as follows;

$$\begin{aligned} s: C \rtimes \mathcal{M}(C) &\longrightarrow \mathcal{M}(C) & t: C \rtimes \mathcal{M}(C) &\longrightarrow \mathcal{M}(C) \\ (x, f) &\longmapsto s(x, f) = f & (x, f) &\longmapsto t(x, f) = M_x \cdot f \end{aligned}$$

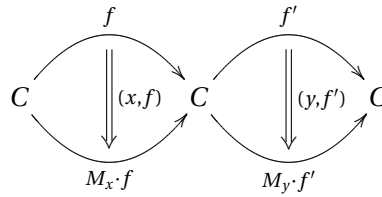
and

$$\begin{aligned} e: \mathcal{M}(C) &\longrightarrow C \rtimes \mathcal{M}(C) \\ f &\longmapsto e(f) = (0, f) \end{aligned}$$

where $M_x \cdot f$ is defined by $(M_x \cdot f)(u) = xu + f(u)$ for $u \in C$.

There are two ways of composing 2-morphisms: vertical and horizontal composition. Now we define this compositions.

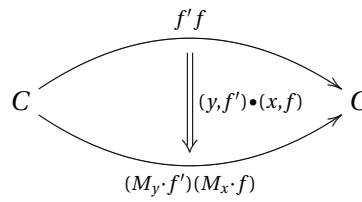
For $(x, f), (y, f') \in C \rtimes \mathcal{M}(C)$



the horizontal composition is defined by

$$(x, f) \bullet (y, f') = (f'(x) + f(y) + xy, f'f),$$

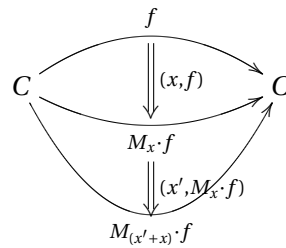
thus we have



and

$$\begin{aligned} t(f'(x) + f(y) + xy, f'f) &= M_{f'(x)+f(y)+xy} \cdot f'f \\ &= (M_y \cdot f')(M_x \cdot f) \end{aligned}$$

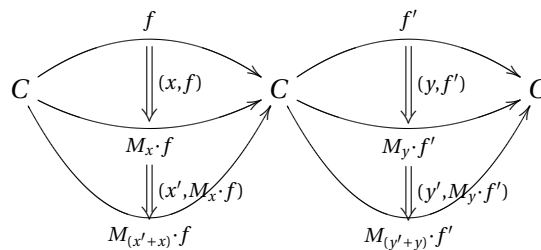
The vertical composition is defined by



$$(x, f) \circ (x', M_x \cdot f) = (x' + x, f)$$

for $(x, f), (x', M_x \cdot f) \in C \rtimes \mathcal{M}(C)$ with $t(x, f) = s(x', M_x \cdot f) = M_x \cdot f$.

It remains to satisfy the interchange law, i.e.



$$\begin{aligned} [(x, f) \circ (x', M_x \cdot f)] \bullet [(y, f') \circ (y', M_y \cdot f')] &= [(x, f) \bullet (y, f')] \\ &\quad \circ [(x', M_x \cdot f) \bullet (y', M_y \cdot f')]. \end{aligned}$$

Evaluating the two sides separately, we get

$$\begin{aligned} \text{LHS} &= (x' + x, f) \bullet (y' + y, f') \\ &= (f'(x' + x) + f(y' + y) + (x' + x)(y' + y), f'f) \\ &= (f'(x') + f'(x) + f(y') + f(y) + x'y' + x'y + xy', f'f) \end{aligned}$$

and

$$\begin{aligned} \text{RHS} &= (f'(x) + f(y) + xy, f'f) \circ ((M_y \cdot f')(x') \\ &\quad + (M_x \cdot f)(y') + x'y', (M_y \cdot f')(M_x \cdot f)) \\ &= (f'(x) + f(y) + xy + (M_y \cdot f')(x') + (M_x \cdot f)(y') + x'y', f'f) \\ &= (f'(x) + f(y) + xy + yx' + f'(x') + xy' + f(y') + x'y', f'f) \end{aligned}$$

LHS and RHS are equal, thus interchange law is satisfied. Therefore we get a 2-algebra consists of the R -algebra C as single object and the R -algebra A_0 of 1-morphisms and the R -algebra A_1 of 2-morphisms.

4. Crossed modules and 2-algebras

Crossed modules have been used widely and in various contexts since their definition by Whitehead [11] in his investigations of the algebraic structure of relative homotopy groups. We recalled the definition of crossed modules of commutative algebras given by Porter [10].

Let R be a k -algebra with identity. A pre-crossed module of commutative algebras is an R -algebra C together with a commutative action of R on C and a morphism

$$\partial : C \longrightarrow R$$

such that for all $c \in C, r \in R$

$$\text{CM1) } \partial(r \blacktriangleright c) = r\partial c.$$

This is a crossed R -module if in addition for all $c, c' \in C$

$$\text{CM2) } \partial c \blacktriangleright c' = cc'.$$

The last condition is called the Peiffer identity. We denote such a crossed module by (C, R, ∂) .

A morphism of crossed modules from (C, R, ∂) to (C', R', ∂') is a pair of k -algebra morphisms $\phi : C \longrightarrow C', \psi : R \longrightarrow R'$ such that

$$\partial' \phi = \psi \partial \quad \text{and} \quad \phi(r \blacktriangleright c) = \psi(r) \blacktriangleright \phi(c).$$

Thus we get a category \mathbf{XMod}_k of crossed modules (for fixed k).

Examples of Crossed Modules

1. Any ideal I in R gives an inclusion map, $inc : I \longrightarrow R$ which is a crossed module. Conversely given an arbitrary R -module $\partial : C \longrightarrow R$ one easily sees that the Peiffer identity implies that ∂C is an ideal in R .
2. Any R -module M can be considered as an R -algebra with zero multiplication and hence the zero mor-

phism $0 : M \rightarrow R$ sending everything in M to the zero element of R is a crossed module. Conversely: If (C, R, ∂) is a crossed module, $\partial(C)$ acts trivially on $\ker \partial$, hence $\ker \partial$ has a natural $R/\partial(C)$ -module structure. As these two examples suggest, general crossed modules lie between the two extremes of ideal and modules. Both aspects are important.

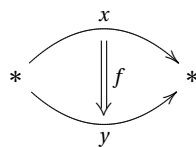
3. Let be $\mathcal{M}(C)$ multiplication algebra. Then $(C, \mathcal{M}(C), \mu)$ is multiplication crossed module. $\mu : C \rightarrow \mathcal{M}(C)$ is defined by $\mu(r) = \delta_r$ with $\delta_r(r') = rr'$ for all $r, r' \in C$, where δ is multiplier $\delta : C \rightarrow C$ such that for all $r, r' \in C$, $\delta(rr') = \delta(r)r'$. Also $\mathcal{M}(C)$ acts on C by $\delta \blacktriangleright r = \delta(r)$. (See [1] for details).

In [10] Porter states that there is an equivalence of categories between the category of internal categories in the category of k -algebras and the category of crossed modules of commutative k -algebras. In the following theorem, we will give a categorical presentation of this equivalence.

Theorem 4.1. The category of crossed modules \mathbf{XMod}_k is equivalent to that of 2-algebras, **2Alg**.

Proof.

Let $A = (A_0, A_1, s, t, e, \circ, \blacktriangleright)$ be a 2-algebra consisting of a single object say $*$ and an algebra A_0 of 1-morphisms and an algebra A_1 of 2-morphisms. For $x, y \in A_0$ and $f : x \rightarrow y \in A_1$, we get the following diagram



We define s, t morphisms $s : A_1 \rightarrow A_0, s(f) = x, t : A_1 \rightarrow A_0, t(f) = y$ and e morphism $e : A_0 \rightarrow A_1$ for $x \in A_0, e(x) : x \rightarrow x \in A_1$.

The s, t and e morphisms are algebra morphisms and we have

$$\begin{aligned}
 se(x) &= s(e(x)) = x = Id_{A_0}(x) \\
 te(x) &= t(e(x)) = x = Id_{A_0}(x)
 \end{aligned}$$

We define

$$\text{Ker } s = H = \{f \in A_1 \mid s(f) = Id_{A_0}\} \subseteq A_1$$

and $\partial = t|_H$ algebra homomorphism by $\partial : H \rightarrow A_0, \partial(h) = t(h)$. We have semidirect product $\text{Ker } s \rtimes A_0 = \{(h, x) \mid h \in \text{Ker } s, x \in A_0\}$ with multiplication $(h, x) \bullet (h', x') = (x \blacktriangleright h' + x' \blacktriangleright h + h' \bullet h, x \bullet x')$ where action of A_0 on $\text{Ker } s$ is defined by $x \blacktriangleright h = e(x) \bullet h$. For each $f \in A_1$, we can write $f = n + e(x)$ where $n = f - es(f) \in \text{Ker } s$ and $x = s(f)$. Suppose $f' = n' + e(x')$. Then

$$\begin{aligned}
 f \bullet f' &= (n + e(x)) \bullet (n' + e(x')) \\
 &= n \bullet n' + n \bullet e(x') + e(x) \bullet n' + e(x) \bullet e(x') \\
 &= e(x') \bullet n + e(x) \bullet n' + n \bullet n' + e(x \bullet x') \\
 &= x' \blacktriangleright n + x \blacktriangleright n' + n \bullet n' + e(x \bullet x').
 \end{aligned}$$

There is a map

$$\begin{aligned}
 \phi : \quad A_1 &\longrightarrow \text{Ker } s \rtimes A_0 \\
 n + e(x) &\longmapsto \phi(n + e(x)) = (n, x).
 \end{aligned}$$

Now

$$\begin{aligned}
 \phi(f \bullet f') &= \phi(x' \blacktriangleright n + x \blacktriangleright n' + n \bullet n' + e(x \bullet x')) \\
 &= (x' \blacktriangleright n + x \blacktriangleright n' + n \bullet n', x \bullet x') \\
 &= (n, x) \bullet (n', x') \\
 &= \phi(f) \bullet \phi(f')
 \end{aligned}$$

so ϕ is a homomorphism. Also, there is an obvious inverse

$$\begin{aligned}
 \phi^{-1}: \text{Kers} \rtimes A_0 &\longrightarrow A_1 \\
 (n, x) &\longmapsto \phi^{-1}(n, x) = n + e(x)
 \end{aligned}$$

which is also a homomorphism. Hence ϕ is an isomorphism and we have established that $\text{Ker } s \rtimes A_0 \cong A_1$.

Since A is a 2-algebra and $\text{Ker } s \rtimes A_0 \cong A_1$, we can define algebra morphisms

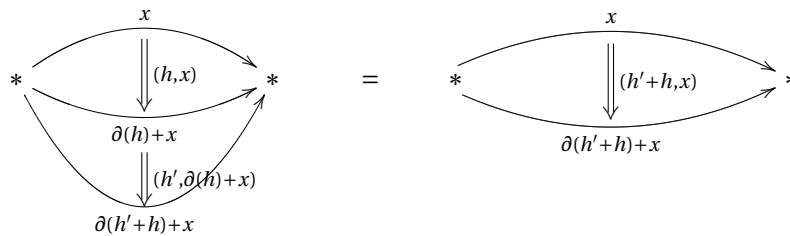
$$\begin{aligned}
 s: \text{Kers} \rtimes A_0 &\longrightarrow A_0 & t: \text{Kers} \rtimes A_0 &\longrightarrow A_0 \\
 (h, x) &\longmapsto s(h, x) = x & (h, x) &\longmapsto t(h, x) = \partial(h) + x
 \end{aligned}$$

and

$$\begin{aligned}
 e: A_0 &\longrightarrow \text{Kers} \rtimes A_0 \\
 x &\longmapsto e(x) = (0, x)
 \end{aligned}$$

and for $t(h, x) = s(h', \partial(h) + x) = \partial(h) + x$ we define

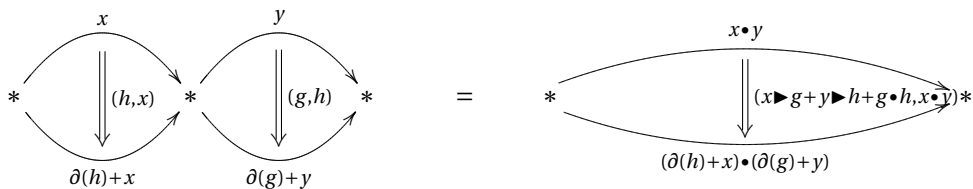
$$\begin{aligned}
 \circ: \text{Kers} \rtimes A_0 \times_s \text{Kers} \rtimes A_0 &\longrightarrow \text{Kers} \rtimes A_0 \\
 ((h, x), (h', \partial(h) + x)) &\longmapsto (h' + h, x)
 \end{aligned}$$



which is vertical composition;

$$(h, x) \circ (h', \partial(h) + x) = (h' + h, x).$$

For $(h, x), (g, y) \in \text{Kers} \rtimes A_0$, horizontal composition is defined by



$$\begin{aligned}
 (h, x) \bullet (g, y) &= (x \blacktriangleright g + y \blacktriangleright h + g \bullet h, x \bullet y) \\
 &= (e(x) \bullet g + e(y) \bullet h + g \bullet h, x \bullet y).
 \end{aligned}$$

Thus we have

CM1)

$$\begin{aligned}\partial(x \blacktriangleright h) &= \partial(e(x) \bullet h) \\ &= \partial(e(x)) \bullet \partial(h) \\ &= (te)(x) \bullet \partial(h) \\ &= x \bullet \partial(h).\end{aligned}$$

Also by interchange law we have

$$\begin{aligned}[(h, x) \bullet (g, y)] \circ [(h', \partial(h) + x) \bullet (g', \partial(g) + y)] &= [(h, x) \circ (h', \partial(h) + x)] \\ &\bullet [(g, y) \circ (g', \partial(g) + y)].\end{aligned}$$

Therefore, evaluating the two sides of this equation gives:

$$\begin{aligned}LHS &= (x \blacktriangleright g + y \blacktriangleright h + g \bullet h, x \bullet y) \\ &\quad \circ ((\partial(h) + x) \blacktriangleright g' + (\partial(g) + y) \blacktriangleright h' + g' \bullet h', (\partial(h) + x) \bullet (\partial(g) + y)) \\ &= ((\partial(h) + x) \blacktriangleright g' + (\partial(g) + y) \blacktriangleright h' + g' \bullet h' + x \blacktriangleright g + y \blacktriangleright h + g \bullet h, x \bullet y) \\ &= (\partial(h) \blacktriangleright g' + e(x) \bullet g' + \partial(g) \blacktriangleright h' \\ &\quad + e(y) \bullet h' + g' \bullet h' + e(x) \bullet g + e(y) \bullet h + g \bullet h, x \bullet y) \\ RHS &= (h' + h, x) \bullet (g' + g, y) \\ &= (x \blacktriangleright (g' + g) + y \blacktriangleright (h' + h) + (g' + g) \bullet (h' + h), x \bullet y) \\ &= (e(x) \bullet g' + e(x) \bullet g + e(y) \bullet h' + e(y) \bullet h + g' \bullet h' + g' \bullet h + g \bullet h' + g \bullet h, x \bullet y).\end{aligned}$$

Since the two sides are equal, we know that their first components must be equal, so we have

$$\partial(h) \blacktriangleright g' + \partial(g) \blacktriangleright h' = h \bullet g' + g \bullet h'$$

and

$$\begin{aligned}h \bullet g' + g \bullet h' &= \partial(h) \blacktriangleright g' + \partial(g) \blacktriangleright h' \\ &= \partial(h + g) \blacktriangleright (g' + h') - \partial(h) \blacktriangleright h' - \partial(g) \blacktriangleright g' \\ &= \partial(h + g) \blacktriangleright (g' + h') - (h \bullet h' + g \bullet g'),\end{aligned}$$

thus

$$\begin{aligned}\partial(h + g) \blacktriangleright (g' + h') &= h \bullet g' + g \bullet h' + (h \bullet h' + g \bullet g') \\ &= (h + g) \bullet (h' + g')\end{aligned}$$

and writing $(h + g) = l, (h' + g') = l' \in Kers$, we get

$$\partial(l) \blacktriangleright l' = l \bullet l'$$

which is the Peiffer identity as required. Hence $(Kers, A_0, \partial)$ is a crossed module.

Let $\mathcal{A} = (A_0, A_1, s, t, e, \circ, \bullet)$ and $\mathcal{A}' = (A'_0, A'_1, s', t', e', \circ', \bullet')$ be 2-algebras and $F = (F_0, F_1) : \mathcal{A} \longrightarrow \mathcal{A}'$ be a 2-algebra morphism. Then $F_0 : A_0 \longrightarrow A'_0$ and $F_1 : A_1 \longrightarrow A'_1$ are the k -algebra morphisms. We define $f_1 =$

$F_1|_{Kers} : Kers \rightarrow Kers'$ and $f_0 = F_0 : A_0 \rightarrow A_0'$. For all $a \in Kers$ and $x \in A_0$,

$$\begin{aligned} f_0 \partial(a) &= F_0 t(a) \\ &= t' F_1(a) \\ &= \partial' f_1(a) \end{aligned}$$

and

$$\begin{aligned} f_1(x \blacktriangleright a) &= F_1(e(x)a) \\ &= F_1(e(x))F_1(a) \\ &= e' F_0(x)F_1(a) \\ &= e' f_0(x)f_1(a) \\ &= f_0(x) \blacktriangleright f_1(a). \end{aligned}$$

Thus (f_1, f_0) map is a crossed module morphism $(Kers, A_0, \partial) \rightarrow (Kers', A_0', \partial')$. So we have a functor

$$\Gamma : \mathbf{2Alg} \rightarrow \mathbf{XMod}_k.$$

Conversely, let (G, C, ∂) be a crossed module of algebras. Therefore there is an algebra morphism $\partial : G \rightarrow C$ and an action of C on G such that

CM1) $\partial(x \blacktriangleright g) = x\partial(g)$,

CM2) $\partial(g) \blacktriangleright g' = gg'$.

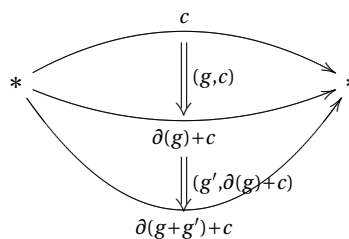
Since C acts on G , we can form the semidirect product $G \rtimes C$ as defined by

$$G \rtimes C = \{(g, c) \mid g \in G, c \in C\}$$

with multiplication

$$(g, c)(g', c') = (c \blacktriangleright g' + c' \blacktriangleright g + g'g, cc')$$

and define maps $s, t : G \rtimes C \rightarrow C$ and $e : C \rightarrow G \rtimes C$ by $s(g, c) = c$, $t(g, c) = \partial(g) + c$ and $e(c) = (0, c)$. These maps are clearly algebra morphisms.



For $t(g, c) = s(g', \partial(g) + c) = \partial(g) + c$, we define composition

$$\begin{aligned} \circ : (G \rtimes C)_t \times_s (G \rtimes C) &\rightarrow (G \rtimes C) \\ (g, c), (g', \partial(g) + c) &\mapsto (g + g', c), \end{aligned}$$

for $(g, c), (h, d) \in G \rtimes C$ and $(g, c), (g', \partial(g) + c) \in G \rtimes C$, following equations give horizontal and vertical composition respectively.

$$(g, c) \bullet (h, d) = (c \blacktriangleright h + d \blacktriangleright g + gh, cd)$$

$$(g, c) \circ (g', \partial(g) + c) = (g + g', c)$$

Finally, since it must be that \circ is an algebra morphism and by the crossed module conditions, interchange law is satisfied. Therefore we have constructed a 2-algebra $\mathcal{A} = (C, G \rtimes C, s, t, e, \circ, \bullet)$ consists of the single object say $*$ and the k -algebra C of 1-morphisms and the k -algebra $G \rtimes C$ of 2-morphisms. Let (G, C, ∂) and (G', C', ∂') be crossed modules and $f = (f_1, f_0) : (G, C, \partial) \rightarrow (G', C', \partial')$ be a crossed module morphism. We define

$$\begin{aligned} F_1 : G \rtimes C &\longrightarrow G' \rtimes C' \\ (g, c) &\longmapsto F_1(g, c) = (f_1(g), f_0(c)) \end{aligned}$$

and

$$\begin{aligned} F_0 : C &\longrightarrow C' \\ c &\longmapsto F_0(c) = f_0(c). \end{aligned}$$

Then

$$\begin{aligned} s' F_1(g, c) &= s'(f_1(g), f_0(c)) \\ &= f_0(c) \\ &= F_0(c) \\ &= F_0 s(g, c), \end{aligned}$$

$$\begin{aligned} t' F_1(g, c) &= t'(f_1(g), f_0(c)) \\ &= \partial' f_1(g) + f_0(c) \\ &= f_0 \partial(g) + f_0(c) \\ &= F_0(\partial(g) + c) \\ &= F_0 t(g, c), \end{aligned}$$

$$\begin{aligned} e' F_0(c) &= (0, f_0(c)) \\ &= F_1(0, c) \\ &= F_1 e(c), \end{aligned}$$

$$\begin{aligned} F_1((g, c) \circ (g', c')) &= F_1(g + g', c) \\ &= (f_1(g + g'), f_0(c)) \\ &= (f_1(g) + f_1(g'), f_0(c)) \\ &= (f_1(g), f_0(c)) \circ (f_1(g'), f_0(c')) \\ &= F_1(g, c) \circ F_1(g', c'), \end{aligned}$$

$$\begin{aligned} F_1((g, c) \bullet (h, d)) &= F_1(c \blacktriangleright h + d \blacktriangleright g + gh, cd) \\ &= (f_1(c \blacktriangleright h) + f_1(d \blacktriangleright g) + f_1(gh), f_0(cd)) \\ &= (f_0(c) \blacktriangleright f_1(h) + f_0(d) \blacktriangleright f_1(g) + f_1(g) f_1(h), f_0(c) f_0(d)) \\ &= (f_1(g), f_0(c)) \bullet (f_1(h), f_0(d)) \\ &= F_1(g, c) \bullet F_1(h, d) \end{aligned}$$

for all $(g, c) \in G \rtimes C$ and $c \in C$. Therefore $F = (F_1, F_0)$ is a 2-algebra morphism from $(C, G \rtimes C, s, t, e, \circ, \bullet)$ to $(C', G' \rtimes C', s', t', e', \circ', \bullet')$. Thus we get a functor

$$\Psi : \mathbf{XMod}_k \longrightarrow \mathbf{2Alg}.$$

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