



---

## Bivariate Variations of Fibonacci and Narayana Sequences and Universal Codes

Çağla Çelemoğlu<sup>1</sup> 

### Article Info

Received: 10 Nov 2022  
Accepted: 23 Dec 2022  
Published: 31 Dec 2022  
doi:10.53570/jnt.1202341  
Research Article

**Abstract** — In this study, we worked on the third-order bivariate variant of the Fibonacci universal code and the second-order bivariate variant of the Narayana universal code, depending on two negative integer variables  $u$  and  $v$ . We then showed in tables these codes for  $1 \leq k \leq 100$ ,  $u = -1, -2, \dots, -20$ , and  $v = -2, -3, \dots, -21$  ( $u$  and  $v$  are consecutive,  $v < u$ ). Moreover, we obtained some significant results from these tables. Furthermore, we compared the use of these codes in cryptography. Finally, we obtained the third-order bivariate variant of Fibonacci codes is more valuable than the second-order bivariate variant of Narayana codes.

**Keywords** — Fibonacci sequence, Narayana sequence, cryptography, variant Fibonacci code, variant Narayana code

**Mathematics Subject Classification (2020)** – 11B37, 14G50

### 1. Introduction

Number sequences have been popular with scientists for centuries. The most popular of these sequences is the Fibonacci sequence. The Fibonacci sequence,  $\{F_k\}_0^\infty$ , is a series of numbers, starting with the integers 0 and 1, in which the value of any element is computed by taking the summation of the two consecutive numbers. Here, for  $k \geq 2$ ,  $F_k = F_{k-1} + F_{k-2}$  [1]. In 1202, this sequence was introduced by Fibonacci. The Fibonacci numbers in the range of  $-8 \leq k \leq 8$  are as follows:

$$-21, 13, -8, 5, -3, 2, -1, 1, 0, 1, 1, 2, 3, 5, 8, 13, 21$$

A sequence that based on the Fibonacci sequence and is a generalization of the Fibonacci sequence is the Tribonacci sequence. The Tribonacci sequence,  $\{T_k\}_0^\infty$ , is a series of numbers, starting with the integers 0, 1, and 1, in which the value of any element is computed by taking the summation of the preceding three terms [2]. If so, for  $k \geq 3$ ,  $T_k = T_{k-1} + T_{k-2} + T_{k-3}$  [1]. In 1914, this number series was presented by Agronomof. The Tribonacci numbers in the range of  $-8 \leq k \leq 8$  are as follows:

$$4, 1, -3, 2, 0, -1, 1, 0, 0, 1, 1, 2, 4, 7, 13, 24, 44$$

Another sequence that based on a similar problem with Fibonacci numbers is the Narayana sequence. The Narayana sequence  $\{N_k\}_0^\infty$ , is a series of numbers, starting with the integers 1, 1, and 1, in which the value of any element is computed by taking the summation of the previous term and term two places before. Then, for  $k \geq 3$ ,  $N_{k+1} = N_k + N_{k-2}$  [1].

---

<sup>1</sup>cagla.ozyilmaz@omu.edu.tr (Corresponding Author)

<sup>1</sup>Department of Mathematics, Faculty of Sciences, Ondokuz Mayıs University, Samsun, Türkiye

In 1356, this number series was presented by Narayana in the treatise named *Gaṇita Kaumudi*. The Narayana numbers in the range of  $-8 \leq k \leq 8$  are as follows:

$$-2, 1, 1, -1, 0, 1, 0, 0, 1, 1, 2, 3, 4, 6, 9, 13$$

There are also different variants of the Fibonacci sequence. The author [3] defined a variation of the Fibonacci sequence that is the more general of Gopala Hemachandra (GH) sequence [4]:

$$\{u, v, u + v, u + 2v, 2u + 3v, 3u + 5v\}$$

where  $u, v$  are integers and a second order variant Fibonacci sequence, as the Gopala Hemachandra sequence above such that  $v = 1 - u$ .

In addition, different universal codes have been defined in the literature. Codes that are obtained by using number sequences are involving to them. The Fibonacci universal code, GH code (or variant Fibonacci code) and the Narayana code are examples of these can be given [3-9]. The most used of these codes is the Fibonacci code [8]. Fibonacci code is a universal code that encodes positive integers with binary representations according to Zeckendorf's Theorem [9]. To find the Fibonacci code for any positive integer, the following algorithm must hold.

One Fibonacci representation or code can be obtained for each positive integer  $C$  with a binary string of length  $t$ ,  $g_1 g_2 \dots g_{t-1} g_t$ , such that  $C = \sum_{i=1}^t g_i F_i^{(m)}$ . The representation is unique if and only if it is used the following algorithm: When it is given the integer  $C$ , it is selected the largest Fibonacci number  $F_k^{(2)}$  equal to  $C$  or smaller to  $C$ ; after that, it is continued repeating with  $C - F_k^{(2)}$  [8]. Finally, to find the Fibonacci code for any positive integer, it is added 1-bit to the Fibonacci representation of the positive integer. For example,  $31 = 2 + 8 + 21 = F_2^{(2)} + F_5^{(2)} + F_7^{(2)}$ . Hence, its Fibonacci representation is 0100101 and its Fibonacci code is 01001011.

According to the above algorithm, there is no contiguous 1-bit in the binary representation. If we apply this algorithm to higher orders, it is seen that the same operations are carried out. But in the higher orders, it is added 1,  $m - 1$  bits to the  $m^{\text{th}}$  order Fibonacci representation of  $k$  to build the  $m^{\text{th}}$  order Fibonacci code. Therefore, there are no successive of  $m$  bits [8].

Narayana and GH universal code are two generalizations of Fibonacci universal code. Thus, these codes also encode positive integers with binary representations according to above algorithm. That is, these universal codes and the variants of these codes can be obtained by using the same rule used to generate the standard Fibonacci code [3].

Moreover, cryptographic applications can be made using these codes or representations. There are many works on these codes in literature (see for example [5-10]).

This paper research the second-order bivariate variant of Narayana codes and the third-order bivariate GH codes and the properties of these codes. Furthermore, it also gives some results regarding the use of these codes in cryptography.

## 2. Preliminaries

In this section, the basic definitions and theorems that used in this study will be given.

**Definition 2.1.** [8] The  $m^{\text{th}}$  order Fibonacci numbers, represented by  $F_k^{(m)}$ , are described with iteration relation as follows:  $F_k^{(m)} = F_{k-1}^{(m)} + F_{k-2}^{(m)} + \dots + F_{k-m}^{(m)}$ , for  $k > 0$ , and the boundary conditions  $F_0^{(m)} = 1$  and  $F_l^{(m)} = 0$  ( $-m < l < 0$ ).

Thus, in 2007, a second-order variation of the Fibonacci code was presented for  $m = 2$  as follows:

**Definition 2.2.** [3] The second-order variant Fibonacci sequences,  $GH_u^{(2)}(k)$  is described with the sequences  $\{u, v, u + v, u + 2v, 2u + 3v, 3u + 5v\}$  where  $v = 1 - u$ , that is ,  $GH_u^{(2)}(1) = u$ ;  $GH_u^{(2)}(2) = 1 - u$ ; and for  $k \geq 3$ ,  $GH_u^{(2)}(k) = GH_u^{(2)}(k - 1) + GH_u^{(2)}(k - 2)$ .

Afterwards, in 2010, the authors [9] studied on the second-order variation Fibonacci code and obtained some results about the second-order variation Fibonacci codes of some positive integers. Then, in 2015, a third-order variation of the Fibonacci code was presented for  $m = 3$  as follows:

**Definition 2.3.** [10] The third-order variant Fibonacci sequences,  $GH_u^{(3)}(k)$  is described with the sequences  $\{u, v, u + v, 2u + 2v, 3u + 4v, 6u + 7v\}$  where  $v = 1 - u$ , that is ,  $GH_u^{(3)}(1) = u$ ,  $GH_u^{(3)}(2) = 1 - u$ ,  $GH_u^{(3)}(3) = 1$ , and for  $k \geq 4$ ,  $GH_u^{(3)}(k) = GH_u^{(3)}(k - 1) + GH_u^{(3)}(k - 2) + GH_u^{(3)}(k - 3)$ .

In this situation, the second and the third order variant Fibonacci (or GH) codes can be defined above definitions. And we know that these universal codes can be obtained by using the same rule used to generate the standard Fibonacci code [3]. Thus, it is obvious that as the value of  $u$  changes, a different sequence is obtained. Daykin [11] proved that only the Fibonacci sequence forms a unique Fibonacci code for all positive integers. Moreover, some integers have many GH codes, while others have no GH code. For instance, if  $VF_{-5}^{(2)}(k) = \{-5, 6, 1, 7, 8, 15, 23, 38, \dots\}$ , then there is no GH code for  $k = 20$ . Similarly, if  $VF_{-2}^{(2)}(k) = \{-2, 3, 1, 4, 5, 9, 14, 23, \dots\}$ , then there are two GH codes for  $k = 13$ . Because

$$13 = 4 + 9 = VF_{-2}^{(2)}(4) + VF_{-2}^{(2)}(6) = -2 + 1 + 14 = VF_{-2}^{(2)}(1) + VF_{-2}^{(2)}(3) + VF_{-2}^{(2)}(7)$$

these codes are 0001011 and 10100011. Furthermore, if  $VF_{-11}^{(3)}(k) = \{-11, 12, 1, 2, 15, 18, 35, 68, \dots\}$ , then there is no GH code for  $k = 11$  [10], and if  $VF_{-5}^{(3)}(k) = \{-5, 6, 1, 2, 9, 12, 23, 44, \dots\}$ , then there are two GH codes for integer  $k = 9$ . Because  $9 = VF_{-5}^{(3)}(5) = -5 + 2 + 12 = VF_{-5}^{(3)}(1) + VF_{-5}^{(3)}(4) + VF_{-5}^{(3)}(6)$  these codes are 0000111 and 10010111. Moreover, in 2016, it was obtained and proved the following theorem by Basu et al [12]. According to following theorem, it is obtained that the relationship between the second-order variant Fibonacci sequence and the Fibonacci sequence.

**Theorem 2.4.** [12] Let  $VF_u(k)$  is the second-order variant Fibonacci sequence and  $F(k)$  is the Fibonacci sequence. In this case, for  $k \geq 1$ ,  $VF_u(k) = F(k - 2) - uF(k - 4)$  where  $k$  is an integer.

**Definition 2.5.** [5] A second-order variant of the Narayana sequence,  $VN_u(k)$ , is described the sequence  $\{u, 3 - u, 1 - u, 1, 4 - u, 5 - 2u, 6 - 2u, 10 - 3u, \dots\}$  with initial conditions  $VN_u(0) = u$ ;  $VN_u(1) = 3 - u$ ;  $VN_u(2) = 1 - u$ ; and for  $k \geq 3$ ,  $VN_u(k) = VN_u(k - 1) + VN_u(k - 3)$ .

Similarly, the second order variant Narayana codes can be defined above definition. Furthermore, we know that these universal codes can be obtained by using the same algorithm used to generate the standard Fibonacci code [3]. For instance, there is no variant Narayana code for 2 for  $VN_{-3}(k) = \{-3, 6, 4, 1, 7, 11, 12, 19\}$ . Moreover, we obtain that there are two variant Narayana codes of 12 for  $VN_{-1}(k) = \{-1, 4, 2, 1, 5, 7, 8, 13, 20, \dots\}$ . Because

$$12 = 4 + 8 = VN_{-1}(2) + VN_{-1}(7) = -1 + 13 = VN_{-5}^{(3)}(1) + VN_{-5}^{(3)}(8)$$

these variant Narayana codes of 12 are 01000011 and 100000011.

### 3. The Research Findings and Discussion

In this study, firstly, we obtained and demonstrated the following theorem, which is the relation between the third-order variant Fibonacci sequence and the Tribonacci sequence.

**Theorem 3.1.** Let  $VF_u^{(3)}(k)$  is the third-order variant Fibonacci sequence and  $T(k)$  is the Tribonacci sequence. In this case, for  $k \geq 1$ ,

$$VF_u^{(3)}(k) = T(k - 1) - uT(k - 4)$$

where  $k$  is an integer.

**Proof.** We obtain the following set from the third-order variant of the Fibonacci sequence.

$$VF_u^{(3)}(k) = \{u, 1 - u, 1, 2, 4 - u, 7 - u, 13 - 2u, 24 - 4u, \dots\}$$

$$VF_u^{(3)}(1) = u = 0 - u(-1) = T(-1) - uT(-3)$$

$$VF_u^{(3)}(2) = 1 - u = 1 - u(1) = T(1) - uT(-2)$$

$$VF_u^{(3)}(3) = 1 = 1 - u(0) = T(2) - uT(-1)$$

The results are correct for  $k = 1, k = 2$ , and  $k = 3$ , as seen above. Suppose that the results are correct for  $k = 1, 2, \dots, n$ . In that case,  $VF_u^{(3)}(n - 2) = T(n - 3) - uT(n - 6)$ ,  $VF_u^{(3)}(n - 1) = T(n - 2) - uT(n - 5)$ , and  $VF_u^{(3)}(n) = T(n - 1) - uT(n - 4)$ . Since

$$\begin{aligned} VF_u^{(3)}(n + 1) &= T(n - 1) - uT(n - 4) + T(n - 2) - uT(n - 5) + T(n - 3) - uT(n - 6) \\ &= T(n - 1) + T(n - 2) + T(n - 3) - uT(n - 4) + T(n - 5) + T(n - 6) \\ &= T(n) - uT(n - 3) = T(n + 1 - 1) - uT(n + 1 - 4) \end{aligned}$$

this equation is true for  $n + 1$ . Therefore, by the induction,  $VF_u^{(3)}(k) = T(k - 1) - uT(k - 4), k \geq 1$ .

### 3.1. The Third Order Bivariate Gopala Hemachandra Sequences and Codes

We will describe a new variation depending on two negative variables of the GH sequence for  $m = 3$  as follows:

**Definition 3.1.1.** The third-order bivariate GH sequence  $GH_{(u,v)}^{(3)}(k)$  is described the sequence  $\{u, v, s, u + v + s, u + 2v + 2s, 2u + 3v + 4s, 4u + 6v + 7s, \dots\}$  where  $s = 1 - u - v$ , with initial conditions  $GH_{(u,v)}^{(3)}(1) = u, GH_{(u,v)}^{(3)}(2) = v, GH_{(u,v)}^{(3)}(3) = 1 - u - v$  and for  $k \geq 4, GH_{(u,v)}^{(3)}(k) = GH_{(u,v)}^{(3)}(k - 1) + GH_{(u,v)}^{(3)}(k - 2) + GH_{(u,v)}^{(3)}(k - 3)$ .

For instance, from the above definition for  $u = -2, v = -3$ , we have  $\{-2, -3, 6, 1, 4, 11, 16, 31, \dots\}$ . It is obviously seen that different sequences will be obtained for different values of  $u$  and  $v$ . With these bivariate variants of GH sequences, we may compose a new universal source code, which we named the bivariate GH code or bivariate variant of Fibonacci universal code. To make, it can be used the same rule used to generate the standard Fibonacci code as in the second and the third order variant Fibonacci code. In addition, we know that GH sequences authorize having more than one Zeckendorf representation of any integer. Similarly, here, we have obtained the third-order bivariate variant of the GH sequences authorizing having more than one Zeckendorf representation of any integer. For example, for  $GH_{(-2,-3)}^{(3)}(k) = \{-2, -3, 6, 1, 4, 11, 16, 31, \dots\}$ ,

$$17 = 6 + 11 = GH_{(-2,-3)}^{(3)}(3) + GH_{(-2,-3)}^{(3)}(6) = 1 + 16 = GH_{(-2,-3)}^{(3)}(4) + GH_{(-2,-3)}^{(3)}(7)$$

Therefore, the third order bivariate GH code of 17 would be both 00100111 and 000100111.

In this study, we investigated for which values  $u$  and  $v$ , the third-order bivariate GH codes precisely exist or for which  $k$  positive integers they do not exist. For instance, we obtained that each positive integer has at least one bivariate GH code for  $GH_{(-2,-6)}^{(3)}(k) = \{-2, -6, 9, 1, 4, 14, 19, 37, \dots\}$ . But there is no bivariate GH code of 5 for  $GH_{(-6,-2)}^{(3)}(k) = \{-6, -2, 9, 1, 8, 18, 27, 53, \dots\}$ .

Basu and Prasad [9] obtained the second-order GH codes  $GH_u^{(2)}(k)$  or undetectable values of the positive integer  $k$  for  $k = 1, 2, \dots, 100$  and for  $u = -2, \dots, -20$ . Besides, Nalh and Özyılmaz [10] obtained the third-order GH codes  $GH_u^{(3)}(k)$  or undetectable values of the positive integer  $k$ , for  $k = 1, 2, \dots, 100$  and for  $u = -2, \dots, -20$ .

This section obtained the third-order bivariate GH codes  $VF_{(u,v)}^{(3)}(k)$  or undetectable values (-) of the positive integer  $k$  for  $k = 1, 2, \dots, 100$  and for  $u = -1, \dots, -20$  and  $v = -2, \dots, -21$  ( $u$  and  $v$  are consecutive,  $v < u$ ) Tables 1 and 2. From Tables 1 and 2, we got the following results for the third-order bivariate GH codes.

- i. For the positive integers  $k = 1, 2, 3$ , the third-order bivariate GH code  $VF_{(u,v)}^{(3)}(k)$  exactly exists, for  $u = -1, -2, \dots, -20$  and  $v = -2, -3, \dots, -21$  ( $u$  and  $v$  are consecutive,  $v < u$ ).
- ii. For  $1 \leq k \leq 100$ , there are at most  $j$  consecutive undetectable (--) values in the third-order bivariate GH code in  $VF_{(-3+j,-(4+j))}^{(3)}(k)$  column in which  $1 \leq j \leq 17$ .
- iii. For  $1 \leq k \leq 100$ , as long as  $j$  raises, the detectable of GH code is reduced in  $VF_{(-3+j,-(4+j))}^{(3)}(k)$  column in which  $1 \leq j \leq 17$ .

**Table 1.** The third order bivariate GH codes ( $u$  and  $v$  are consecutive and  $v < u$ )

$k$	$(u = -1)$ $(v = -2)$	$(u = -2)$ $(v = -3)$	$(u = -3)$ $(v = -4)$	$(u = -4)$ $(v = -5)$	$(u = -5)$ $(v = -6)$	$(u = -6)$ $(v = -7)$	$(u = -7)$ $(v = -8)$	$(u = -8)$ $(v = -9)$	$(u = -9)$ $(v = -10)$	$(u = -10)$ $(v = -11)$
1	000111	000111	000111	000111	000111	000111	000111	000111	000111	000111
2	1000111	1000111	1000111	1000111	1000111	1000111	1000111	1000111	1000111	1000111
3	0000111	1001111	1001111	1001111	1001111	1001111	1001111	1001111	1001111	1001111
4	00111	0000111	01111	-	-	-	-	-	-	-
5	001111	0001111	0000111	01111	-	-	-	-	-	-
6	01000111	00111	0001111	0000111	01111	-	-	-	-	-
7	10000111	0110111	11000111	0001111	0000111	01111	-	-	-	-
8	00000111	01000111	00111	11000111	0001111	0000111	01111	-	-	-
9	00010111	10000111	001111	11010111	11000111	0001111	0000111	01111	-	-
10	010000111	10010111	01000111	00111	11010111	11000111	0001111	0000111	01111	-
11	100000111	00000111	10000111	001111	-	11010111	11000111	0001111	0000111	01111
12	000000111	00010111	10010111	01000111	00111	-	11010111	11000111	0001111	0000111
13	000100111	010000111	0010111	10000111	001111	-	-	11010111	11000111	0001111
14	100010111	100000111	00000111	10010111	01000111	00111	-	-	11010111	11000111
15	000010111	100100111	00010111	110000111	10000111	001111	-	-	-	11010111
16	001000111	000000111	010000111	0010111	10010111	01000111	00111	-	-	-
17	001100111	000100111	100000111	00000111	110000111	10000111	001111	-	-	-
18	101010111	100010111	100100111	00010111	110100111	10010111	01000111	00111	-	-
19	100001111	100110111	00001111	010000111	0010111	110000111	10000111	001111	-	-
20	000001111	000010111	000000111	100000111	00000111	110100111	10010111	01000111	00111	-
21	0100000111	000110111	000100111	100100111	00010111	-	110000111	10000111	001111	-
22	0101000111	001000111	100010111	01100111	010000111	0010111	110100111	10010111	01000111	00111
23	0000000111	001100111	100110111	00001111	100000111	00000111	-	110000111	10000111	001111
24	0001000111	101010111	011000111	000000111	100100111	00010111	-	110100111	10010111	01000111
25	1000100111	100001111	000010111	000100111	110110111	010000111	0010111	-	110000111	10000111

Table 1. Continued

$k$	$(u = -1)$ $(v = -2)$	$(u = -2)$ $(v = -3)$	$(u = -3)$ $(v = -4)$	$(u = -4)$ $(v = -5)$	$(u = -5)$ $(v = -6)$	$(u = -6)$ $(v = -7)$	$(u = -7)$ $(v = -8)$	$(u = -8)$ $(v = -9)$	$(u = -9)$ $(v = -10)$	$(u = -10)$ $(v = -11)$
26	0000100111	100101111	000110111	100010111	01100111	100000111	00000111	-	110100111	10010111
27	0010000111	000001111	00101111	100110111	00001111	100100111	00010111	-	-	110000111
28	0011000111	0100000111	001000111	00110111	000000111	110110111	010000111	0010111	-	110100111
29	0100010111	0101000111	001100111	011000111	000100111	-	100000111	00000111	-	-
30	0101010111	1001000111	101010111	000010111	100010111	01100111	100100111	00010111	-	-
31	0000010111	0101000111	100001111	000110111	100110111	00001111	110110111	010000111	0010111	-
32	0001010111	0001000111	100101111	110001111	00100111	000000111	-	100000111	00000111	-
33	0100001111	1000100111	001010111	00101111	00110111	000100111	-	100100111	00010111	-
34	0101001111	1001100111	000001111	001000111	011000111	100010111	01100111	110110111	010000111	0010111
35	0000001111	0000100111	0100000111	001100111	000010111	100110111	00001111	-	100000111	00000111
36	0001001111	0001100111	0101000111	101010111	000110111	-	000000111	-	100100111	00010111
37	1000101111	0010000111	1001000111	100001111	110001111	00100111	000100111	-	110110111	010000111
38	0000101111	0011000111	011001111	100101111	110101111	00110111	100010111	01100111	-	100000111
39	0010001111	0100010111	0101000111	1101000111	00101111	011000111	100110111	00001111	-	100100111
40	0011001111	1000010111	0001000111	001010111	001000111	000010111	-	000000111	-	110110111
41	01000000111	1001010111	1000100111	000001111	001100111	000110111	-	000100111	-	-
42	01010000111	0000010111	1001100111	0100000111	101010111	110001111	00100111	100010111	01100111	-
43	00000000111	0001010111	0110000111	0101000111	100001111	110101111	00110111	100110111	00001111	-
44	00010000111	0100001111	0000100111	1001000111	100101111	-	011000111	-	000000111	-
45	10001000111	1000001111	0001100111	-	1101000111	00101111	000010111	000010111	000100111	-
46	00001000111	1001001111	1100010111	011001111	-	001000111	000110111	-	100010111	01100111
47	00100000111	0000001111	0010000111	0101000111	001010111	001100111	110001111	00100111	100110111	00001111
48	00110000111	0001001111	0011000111	0001000111	000001111	101010111	110101111	00110111	-	000000111
49	01000100111	1000101111	0100010111	1000100111	0100000111	100001111	-	011000111	-	000100111
50	10000100111	1001101111	1000010111	-	0101000111	100101111	-	000010111	000010111	100010111
51	00000100111	0000101111	1001010111	001001111	1001000111	1101000111	00101111	000110111	-	100110111
52	00010100111	0001101111	1100001111	0110000111	1101100111	-	001000111	110001111	00100111	-
53	10001100111	11000000111	0000010111	0000100111	-	-	001100111	110101111	00110111	-
54	00001100111	0110101111	0001010111	0001100111	011001111	001010111	101010111	-	011000111	-
55	00000010111	01000000111	1000110111	1100010111	0101000111	000001111	100001111	-	000010111	000010111
56	00010010111	10000000111	1000001111	1101010111	0001000111	0100000111	100101111	-	000110111	
57	10001010111	10010000111	1001001111	0010000111	1000100111	0101000111	1101000111	00101111	110001111	00100111
58	00001010111	00000000111	0000110111	0011000111	1001100111	1001000111	-	001000111	110101111	00110111
59	00100010111	00010000111	0000001111	0100010111	-	1101100111	-	001100111	-	011000111
60	00110010111	10001000111	0001001111	1000010111	001001111	-	-	101010111	-	000010111
61	01000110111	01100000111	1000101111	1001010111	0110000111	-	001010111	100001111	-	000110111
62	10000110111	00001000111	1001101111	1100001111	0000100111	011001111	000001111	100101111	-	110001111
63	00000110111	00011000111	0110001111	0010100111	0001100111	0101000111	0100000111	1101000111	00101111	110101111
64	01000001111	00100000111	0000101111	0000010111	1100010111	0001000111	0101000111	-	001000111	-
65	10000001111	00110000111	0001101111	0001010111	1101010111	1000100111	1001000111	-	001100111	-
66	00000001111	10101000111	11000000111	1000110111	-	-	1101100111	-	101010111	-
67	00010001111	10000100111	11010000111	1000001111	0010000111	-	-	-	100001111	-
68	10001001111	10010100111	0110101111	1001001111	0011000111	-	-	001010111	100101111	-
69	00001001111	00000100111	01000000111	0110010111	0100010111	001001111	-	000001111	1101000111	00101111
70	00100001111	00010100111	10000000111	0000110111	1000010111	0110000111	011001111	0100000111	1001100111	001000111
71	00110001111	10001100111	10010000111	0000001111	1001010111	0000100111	0101000111	0101000111	-	001100111
72	01000101111	10000010111	11011000111	0001001111	1100001111	0001100111	0001000111	1001000111	-	101010111

**Table 1.** Continued

$k$	$(u = -1)$ $(v = -2)$	$(u = -2)$ $(v = -3)$	$(u = -3)$ $(v = -4)$	$(u = -4)$ $(v = -5)$	$(u = 5)$ $(v = -6)$	$(u = -6)$ $(v = -7)$	$(u = -7)$ $(v = -8)$	$(u = -8)$ $(v = -9)$	$(u = -9)$ $(v = -10)$	$(u = -10)$ $(v = -11)$
73	01010101111	00001100111	00000000111	1000101111	1101001111	1100010111	1000100111	1101100111	-	100001111
74	000001001111	00000010111	00010000111	1001101111	0010100111	1101010111	1001100111	-	-	100101111
75	110000000111	00010010111	10001000111	0011010111	0000010111	-	-	-	001010111	1101000111
76	010000000111	10001010111	10011000111	0110001111	0001010111	-	-	-	000001111	-
77	100000000111	01100010111	01100000111	0000101111	1000110111	0010000111	-	-	0100000111	-
78	000000000111	00001010111	00001000111	0001101111	1000001111	0011000111	001001111	011001111	0101000111	-
79	000100000111	00011010111	00011000111	11000000111	1001001111	0100010111	0110000111	0101000111	1001000111	-
80	100010000111	00100010111	11000100111	11010000111	1101101111	1000010111	0000100111	0001000111	1101100111	-
81	000010000111	00110010111	00100000111	0010001111	0110010111	1001010111	0001100111	1000100111	-	-
82	001000000111	01000110111	00110000111	0110101111	0000110111	1100001111	1100010111	1001100111	-	001010111
83	001100000111	10000110111	10101000111	01000000111	0000001111	1101001111	1101010111	-	-	000001111
84	010001000111	11000001111	10000100111	10000000111	0001001111	-	-	-	-	0100000111
85	100001000111	00000110111	10010100111	10010000111	1000101111	0010100111	-	-	-	0101000111
86	000001000111	01000001111	00101000111	11011000111	1001101111	0000010111	-	-	011001111	1001000111
87	000101000111	10000001111	00000100111	0010101111	0010010111	0001010111	0010000111	001001111	0101000111	1101100111
88	010000100111	10010001111	00010100111	00000000111	0011010111	1000110111	0011000111	0110000111	0001000111	-
89	100000100111	00000001111	10001100111	00010000111	0110001111	1000001111	0100010111	0000100111	1000100111	-
90	000000100111	00010001111	10000010111	10001000111	0000101111	1001001111	1000010111	0001100111	1001100111	-
91	000100100111	10001001111	10010010111	10011000111	0001101111	-	1001010111	1100010111	-	-
92	100010100111	01100001111	00001100111	-	11000000111	-	1100001111	1101010111	-	-
93	000010100111	00001001111	00000010111	01100000111	11010000111	0110010111	1101001111	-	-	-
94	001000100111	00011001111	00010010111	00001000111	0010110111	0000110111	-	-	-	011001111
95	001100100111	00100001111	10001010111	00011000111	0010001111	0000001111	-	-	-	0101000111
96	010001100111	00110001111	10011010111	11000100111	0110101111	0001001111	0010100111	-	001001111	0001000111
97	100001100111	01000101111	01100010111	11010100111	01000000111	1000101111	0000010111	0010000111	0110000111	1000100111
98	000001100111	01010101111	00001010111	00100000111	01010000111	1001101111	0001010111	0011000111	0000100111	1001100111
99	000101100111	00101001111	00011010111	00110000111	10010000111	-	1000110111	0100010111	0001100111	-
100	100000010111	110000000111	11000110111	10101000111	11011000111	0010010111	1000001111	1000010111	1100010111	-

**Table 2.** The third order bivariate GH codes ( $u$  and  $v$  are consecutive and  $v < u$ )

$k$	$(u = -11)$ $(v = -12)$	$(u = -12)$ $(v = -13)$	$(u = -13)$ $(v = -14)$	$(u = -14)$ $(v = -15)$	$(u = 15)$ $(v = -16)$	$(u = -16)$ $(v = -17)$	$(u = -17)$ $(v = -18)$	$(u = -18)$ $(v = -19)$	$(u = -19)$ $(v = -20)$	$(u = -20)$ $(v = 21)$
1	000111	000111	000111	000111	000111	000111	000111	000111	000111	000111
2	1000111	1000111	1000111	1000111	1000111	1000111	1000111	1000111	1000111	1000111
3	1001111	1001111	1001111	1001111	1001111	1001111	1001111	1001111	1001111	1001111
4	-	-	-	-	-	-	-	-	-	-
5	-	-	-	-	-	-	-	-	-	-
6	-	-	-	-	-	-	-	-	-	-
7	-	-	-	-	-	-	-	-	-	-
8	-	-	-	-	-	-	-	-	-	-
9	-	-	-	-	-	-	-	-	-	-
10	-	-	-	-	-	-	-	-	-	-
11	-	-	-	-	-	-	-	-	-	-
12	01111	-	-	-	-	-	-	-	-	-
13	0000111	01111	-	-	-	-	-	-	-	-
14	0001111	0000111	01111	-	-	-	-	-	-	-

**Table 2.** Continued

$k$	$(u = -11)$	$(u = -12)$	$(u = -13)$	$(u = -14)$	$(u = -15)$	$(u = -16)$	$(u = -17)$	$(u = -18)$	$(u = -19)$	$(u = -20)$
	$(v = -12)$	$(v = -13)$	$(v = -14)$	$(v = -15)$	$(v = -16)$	$(v = -17)$	$(v = -18)$	$(v = -19)$	$(v = -20)$	$(v = -21)$
15	11000111	0001111	0000111	01111	-	-	-	-	-	-
16	11010111	11000111	0001111	0000111	01111	-	-	-	-	-
17	-	11010111	11000111	0001111	0000111	01111	-	-	-	-
18	-	-	11010111	11000111	0001111	0000111	01111	-	-	-
19	-	-	-	11010111	11000111	0001111	0000111	01111	-	-
20	-	-	-	-	11010111	11000111	0001111	0000111	01111	-
21	-	-	-	-	-	11010111	11000111	0001111	0000111	01111
22	-	-	-	-	-	-	11010111	11000111	0001111	0000111
23	-	-	-	-	-	-	-	11010111	11000111	0001111
24	00111	-	-	-	-	-	-	-	11010111	11000111
25	001111	-	-	-	-	-	-	-	-	11010111
26	01000111	00111	-	-	-	-	-	-	-	-
27	10000111	001111	-	-	-	-	-	-	-	-
28	10010111	01000111	00111	-	-	-	-	-	-	-
29	110000111	10000111	001111	-	-	-	-	-	-	-
30	110100111	10010111	01000111	00111	-	-	-	-	-	-
31	-	110000111	10000111	001111	-	-	-	-	-	-
32	-	110100111	10010111	01000111	00111	-	-	-	-	-
33	-	-	110000111	10000111	001111	-	-	-	-	-
34	-	-	110100111	10010111	01000111	00111	-	-	-	-
35	-	-	-	110000111	10000111	001111	-	-	-	-
36	-	-	-	110100111	10010111	01000111	00111	-	-	-
37	0010111	-	-	-	110000111	10000111	001111	-	-	-
38	00000111	-	-	-	110100111	10010111	01000111	00111	-	-
39	00010111	-	-	-	-	110000111	10000111	001111	-	-
40	010000111	0010111	-	-	-	110100111	10010111	01000111	00111	-
41	100000111	00000111	-	-	-	-	110000111	10000111	001111	-
42	100100111	00010111	-	-	-	-	110100111	10010111	01000111	00111
43	110110111	010000111	0010111	-	-	-	-	110000111	10000111	001111
44	-	100000111	00000111	-	-	-	-	110100111	10010111	01000111
45	-	100100111	00010111	-	-	-	-	-	110000111	10000111
46	-	110110111	010000111	0010111	-	-	-	-	110100111	10010111
47	-	-	100000111	00000111	-	-	-	-	-	110000111
48	-	-	100100111	00010111	-	-	-	-	-	110100111
49	-	-	110110111	010000111	0010111	-	-	-	-	-
50	01100111	-	-	100000111	00000111	-	-	-	-	-
51	00001111	-	-	100100111	00010111	-	-	-	-	-
52	000000111	-	-	110110111	010000111	0010111	-	-	-	-
53	000100111	-	-	-	100000111	00000111	-	-	-	-
54	100010111	01100111	-	-	100100111	00010111	-	-	-	-
55	100110111	00001111	-	-	110110111	010000111	0010111	-	-	-
56	-	000000111	-	-	-	100000111	00000111	-	-	-
57	-	000100111	-	-	-	100100111	00010111	-	-	-
58	-	100010111	01100111	-	-	110110111	010000111	0010111	-	-
59	-	100110111	00001111	-	-	-	100000111	00000111	-	-
60	-	-	000000111	-	-	-	100100111	00010111	-	-
61	-	-	000100111	-	-	-	110110111	010000111	0010111	-



**Table 2.** Continued

$k$	$(u = -11)$ $(v = -12)$	$(u = -12)$ $(v = -13)$	$(u = -13)$ $(v = -14)$	$(u = -14)$ $(v = -15)$	$(u = -15)$ $(v = -16)$	$(u = -16)$ $(v = -17)$	$(u = -17)$ $(v = -18)$	$(u = -18)$ $(v = -19)$	$(u = -19)$ $(v = -20)$	$(u = -20)$ $(v = 21)$
62	00100111	–	100010111	01100111	–	–	–	100000111	00000111	–
63	00110111	–	100110111	00001111	–	–	–	100100111	00010111	–
64	011000111	–	–	000000111	–	–	–	110110111	010000111	0010111
65	000010111	–	–	000100111	–	–	–	–	100000111	00000111
66	000110111	–	–	100010111	01100111	–	–	–	100100111	00010111
67	110001111	00100111	–	100110111	00001111	–	–	–	110110111	010000111
68	110101111	00110111	–	–	000000111	–	–	–	–	100000111
69	–	011000111	–	–	000100111	–	–	–	–	100100111
70	–	000010111	–	–	100010111	01100111	–	–	–	110110111
71	–	000110111	–	–	100110111	00001111	–	–	–	–
72	–	110001111	00100111	–	–	000000111	–	–	–	–
73	–	110101111	00110111	–	–	000100111	–	–	–	–
74	–	–	011000111	–	–	100010111	01100111	–	–	–
75	00101111	–	000010111	–	–	100110111	00001111	–	–	–
76	001000111	–	000110111	–	–	–	000000111	–	–	–
77	011010111	–	110001111	00100111	–	–	000100111	–	–	–
78	101010111	–	110101111	00110111	–	–	100010111	01100111	–	–
79	100001111	–	–	011000111	–	–	100110111	00001111	–	–
80	100101111	–	–	000010111	–	–	–	000000111	–	–
81	1101000111	00101111	–	000110111	–	–	–	000100111	–	–
82	–	001000111	–	110001111	00100111	–	–	100010111	01100111	–
83	–	001100111	–	110101111	00110111	–	–	100110111	00001111	–
84	–	101010111	–	–	011000111	–	–	–	000000111	–
85	–	100001111	–	–	000010111	–	–	–	000100111	–
86	–	100101111	–	–	000110111	–	–	–	100010111	01100111
87	–	1101000111	00101111	–	110001111	00100111	–	–	100110111	00001111
88	–	–	001000111	–	110101111	00110111	–	–	–	000000111
89	001010111	–	001100111	–	–	011000111	–	–	–	000100111
90	000001111	–	101010111	–	–	000010111	–	–	–	100010111
91	000101111	–	100001111	–	–	000110111	–	–	–	100110111
92	0101000111	–	100101111	–	–	110001111	00100111	–	–	–
93	1001000111	–	1101000111	00101111	–	110101111	00110111	–	–	–
94	1101100111	–	–	001000111	–	–	011000111	–	–	–
95	–	–	–	001100111	–	–	000010111	–	–	–
96	–	001010111	–	101010111	–	–	000110111	–	–	–
97	–	000001111	–	100001111	–	–	110001111	00100111	–	–
98	–	0100000111	–	100101111	–	–	110101111	00110111	–	–
99	–	1000000111	–	1101000111	00101111	–	–	011000111	–	–
100	–	1001000111	–	–	001000111	–	–	000010111	–	–

In this section, we have also obtained the following theorem, which is the relation between the third-order bivariate variant of the Fibonacci sequence and the Tribonacci sequence.

**Theorem 3.1.2.** Let  $VF_{(u,v)}^{(3)}(k)$  is the third-order bivariate variant of the Fibonacci sequence and  $T(k)$  is the Tribonacci sequence. In this case, for integers  $k \geq 1$ ,

$$VF_{(u,v)}^{(3)}(k) = T(k - 2) - u(T(k - 4) + T(k - 5)) - vT(k - 5)$$

PROOF. We have the following set from the third-order bivariate variant of the Fibonacci sequence.

$$VF_{(u,v)}^{(3)}(k) = \{u, v, 1 - u - v, 1, 2 - u, 4 - 2u - v, 7 - 3u - v, 13 - 6u - 2v, 24 - 11u - 4v, \dots\}$$

$$VF_{(u,v)}^{(3)}(1) = u = 0 - u(0 + (-1)) - v \cdot 0 = T(-1) - u \cdot (T(-3) + T(-4)) - vT(-4)$$

$$VF_{(u,v)}^{(3)}(2) = v = 0 - u((-1) + 1) - v \cdot (-1) = T(0) - u \cdot (T(-2) + T(-3)) - vT(-3)$$

$$VF_{(u,v)}^{(3)}(3) = 1 - u - v = 1 - u(0 + 1) - v \cdot 1 = T(1) - u \cdot (T(-1) + T(-2)) - vT(-2)$$

The results are correct for  $k = 1, k = 2$ , and  $k = 3$ , as seen above. Suppose that the results are correct for  $k = 1, 2, \dots, n$ . In that case,

$$VF_{(u,v)}^{(3)}(n - 2) = T(n - 4) - u(T(n - 6) + T(n - 7)) - vT(n - 7)$$

$$VF_{(u,v)}^{(3)}(n - 1) = T(n - 3) - u(T(n - 5) + T(n - 6)) - vT(n - 6)$$

$$VF_{(u,v)}^{(3)}(n) = T(n - 2) - u(T(n - 4) + T(n - 5)) - vT(n - 5)$$

Let us show that this equation is true for  $n + 1$ .

$$\begin{aligned} VF_{(u,v)}^{(3)}(n + 1) &= T(n - 4) - u(T(n - 6) + T(n - 7)) - vT(n - 7) + T(n - 3) - u(T(n - 5) + T(n - 6)) \\ &\quad - vT(n - 6) + T(n - 2) - u(T(n - 4) + T(n - 5)) - vT(n - 5) \\ &= T(n - 4) + T(n - 3) + T(n - 2) - u(T(n - 4) + T(n - 5) + T(n - 5)) \\ &\quad + T(n - 6) + T(n - 6) + T(n - 7)) - v(T(n - 7) + T(n - 6) + T(n - 5)) \\ &= T(n - 1) - u(T(n - 3) + T(n - 4)) - vT(n - 4) \\ &= T(n + 1 - 2) - u(T(n + 1 - 4) + T(n + 1 - 5)) - vT(n + 1 - 5) \end{aligned}$$

Therefore, by the induction,  $VF_{(u,v)}^{(3)}(k) = T(k - 2) - u(T(k - 4) + T(k - 5)) - vT(k - 5), k \geq 1$ .

### 3.2. The Second-Order Bivariate Variant of Narayana Sequences and Codes

This section describes a new variation depends on two negative variables of the Narayana sequence as follows:

**Definition 3.2.1.** The second-order bivariate variant of the Narayana sequence  $VN_{(u,v)}^{(2)}(k)$  is described the sequence  $\{u, v, s, u + s, v + u + s, v + u + 2s, v + 2u + 3s, \dots\}$  where  $s = 1 - u - v$ , with initial conditions  $VN_{(u,v)}^{(2)}(1) = u; VN_{(u,v)}^{(2)}(2) = v; VN_{(u,v)}^{(2)}(3) = 1 - u - v$ ; and for  $k \geq 4, VN_{(u,v)}^{(2)}(k) = VN_{(u,v)}^{(2)}(k - 1) + VN_{(u,v)}^{(2)}(k - 3)$ .

Here, if  $u = -2$  and  $v = -3$ , then the second-order bivariate variant of the Narayana sequence is  $\{-2, -3, 6, 4, 1, 7, 11, 12, 19, \dots\}$ . It is obviously seen that different sequences will be obtained for different values of  $u$  and  $v$ .

With these bivariate variant of Narayana sequences, we may compose a new universal source code, which we have named the bivariate variant of the Narayana code. To make, it can be used the same rule used to generate the standard Fibonacci code as in the second and the third order bivariate variant of Fibonacci code.

Moreover, as in the variations of Fibonacci sequences, the bivariate variant of Narayana sequences has more than one Zeckendorf representation of any integer, too. For example, the second order bivariate variant of

Narayana code of 17 would be 1000000011 or 00100011 for  $VN_{(-2,-3)}^{(2)}(k) = \{-2, -3, 6, 4, 1, 7, 11, 12, 19, \dots\}$ . Because

$$17 = -2 + 19 = VN_{(-2,-3)}^{(2)}(1) + VN_{(-2,-3)}^{(2)}(9) = 6 + 11 = VN_{(-2,-3)}^{(2)}(3) + VN_{(-2,-3)}^{(2)}(7)$$

This study has also investigated for which values  $u$  and  $v$  the second-order bivariate variant of Narayana codes precisely exist or for which  $k$  positive integers they don't exist. For instance, we have obtained that there is no bivariate variant of the Narayana code of 2 for  $VN_{(-1,-3)}^{(2)}(k) = \{-1, -3, 5, 4, 1, 6, 10, 11, 17\}$ .

Das and Sinha have obtained the second-order variant of Narayana codes  $VN_u^{(2)}(k)$  or undetectable values of the positive integer  $k$  for  $k = 1, 2, \dots, 50$  and for  $u = -1, -2, \dots, -20$  in [13]. This section has obtained the second-order bivariate variant of Narayana codes  $VN_{(u,v)}^{(2)}(k)$  or undetectable values (-) of the positive integer  $k$  for  $k = 1, 2, \dots, 100$  and for  $u = -1, -2, \dots, -20$  and  $v = -2, -3, \dots, -21$  ( $u$  and  $v$  are consecutive,  $v < u$ ) Tables 3 and 4. From Tables 3 and 4, we got the following results for the second-order bivariate variant of Narayana codes.

- i. For the positive integers  $k = 1$ , the second-order bivariate variant of the Narayana code  $VN_{(u,v)}^{(2)}(k)$  exactly exists for  $u = -1, -2, \dots, -20$  and  $v = -2, -3, \dots, -21$  ( $u$  and  $v$  are consecutive,  $v < u$ ).
- ii. For  $1 \leq k \leq 100$ , there are at most  $j$  consecutive undetectable (-) values in the second-order bivariate variant of Narayana code in  $VN_{(-(1+j), -(2+j))}^{(2)}(k)$  column in which  $1 \leq j \leq 19$ .
- iii. For  $1 \leq k \leq 100$ , as long as  $j$  raises, the detectable of Narayana code is reduced in  $VN_{(-(1+j), -(2+j))}^{(2)}(k)$  column in which  $1 \leq j \leq 19$ .

**Table 3.** The second order bivariate variant of Narayana codes ( $u$  and  $v$  are consecutive and  $v < u$ )

$k$	$(u = -1)$ $(v = -2)$	$(u = -2)$ $(v = -3)$	$(u = -3)$ $(v = -4)$	$(u = -4)$ $(v = -5)$	$(u = -5)$ $(v = -6)$	$(u = -6)$ $(v = -7)$	$(u = -7)$ $(v = -8)$	$(u = -8)$ $(v = -9)$	$(u = -9)$ $(v = -10)$	$(u = -10)$ $(v = -11)$
1	000011	000011	000011	000011	000011	000011	000011	000011	000011	000011
2	10011	10011	10011	10011	10011	10011	10011	10011	10011	10011
3	00011	-	-	-	-	-	-	-	-	-
4	0011	00011	-	-	-	-	-	-	-	-
5	0000011	1000011	00011	-	-	-	-	-	-	-
6	01000011	0011	1000011	00011	-	-	-	-	-	-
7	10000011	0000011	-	1000011	00011	-	-	-	-	-
8	00000011	01000011	0011	-	1000011	00011	-	-	-	-
9	000000011	100100011	0000011	-	-	1000011	00011	-	-	-
10	000010011	100000011	01000011	0011	-	-	1000011	00011	-	-
11	100100011	00000011	100100011	0000011	-	-	-	1000011	00011	-
12	0100000011	000000011	100000011	01000011	0011	-	-	-	1000011	00011
13	1000000011	000010011	100010011	1001011	0000011	-	-	-	-	1000011
14	0000000011	100100011	00000011	100000011	01000011	0011	-	-	-	-
15	0000100011	00010011	000000011	100010011	1001011	0000011	-	-	-	-
16	1001000011	0100000011	000010011	-	100000011	01000011	0011	-	-	-
17	0001000011	1000000011	0010011	00000011	100010011	1001011	0000011	-	-	-
18	0010000011	1000100011	-	000000011	-	100000011	01000011	0011	-	-
19	0000010011	0000000011	00010011	000010011	-	100010011	1001011	0000011	-	-
20	01000000011	0000100011	0100000011	100100011	00000011	-	100000011	01000011	0011	-
21	10000000011	1001000011	1000000011	0010011	000000011	-	100010011	1001011	0000011	-
22	00000000011	-	1000100011	-	000010011	-	-	100000011	01000011	0011
23	00001000011	0001000011	001000011	00010011	100100011	00000011	-	100010011	1001011	0000011

**Table 3. Continued**

$k$	$(u = -1)$ $(v = -2)$	$(u = -2)$ $(v = -3)$	$(u = -3)$ $(v = -4)$	$(u = -4)$ $(v = -5)$	$(u = 5)$ $(v = -6)$	$(u = -6)$ $(v = -7)$	$(u = -7)$ $(v = -8)$	$(u = -8)$ $(v = -9)$	$(u = -9)$ $(v = -10)$	$(u = -10)$ $(v = -11)$
24	10010000011	1000010011	0000000011	0100000011	–	000000011	–	–	100000011	01000011
25	00010000011	0010000011	0000100011	1000000011	0010011	000010011	–	–	100010011	1001011
26	00100000011	0000010011	1001000011	1000100011	–	100100011	00000011	–	–	100000011
27	00000100011	0100000011	–	00100011	00010011	–	000000011	–	–	100010011
28	01000010011	1000000011	–	001000011	0100000011	–	000010011	–	–	–
29	010000000011	10001000011	0001000011	0000000011	1000000011	0010011	100100011	00000011	–	–
30	100000000011	0000000011	1000010011	0000100011	1000100011	–	–	000000011	–	–
31	000000000011	00001000011	–	1001000011	–	00010011	–	000010011	–	–
32	000010000011	10010000011	0010000011	–	00100011	0100000011	–	100100011	00000011	–
33	100100000011	–	0000010011	–	00101011	1000000011	0010011	–	000000011	–
34	000100000011	00010000011	01000000011	–	0000000011	1000100011	–	–	000010011	–
35	001000000011	10000100011	10000000011	0001000011	0000100011	–	00010011	–	100100011	00000011
36	000001000011	00100000011	10001000011	1000010011	1001000011	–	0100000011	–	–	000000011
37	010000100011	00000100011	–	–	–	00100011	1000000011	0010011	–	000010011
38	100000100011	01000010011	00000000011	–	–	001000011	1000100011	–	–	100100011
39	000000100011	010000010011	00001000011	0010000011	–	0000000011	–	00010011	–	–
40	000000010011	100000000011	10010000011	0010100011	–	0000100011	–	0100000011	–	–
41	000010010011	100010000011	–	01000000011	0001000011	1001000011	–	1000000011	0010011	–
42	100100010011	000000000011	–	10000000011	1000010011	–	00100011	1000100011	–	–
43	0100000000011	000010000011	00010000011	10001000011	–	–	001000011	–	00010011	–
44	1000000000011	100100000011	10000100011	–	–	–	0000000011	–	0100000011	–
45	0000000000011	00010010011	–	–	–	–	0000100011	–	1000000011	0010011
46	0000100000011	000100000011	00100000011	00000000011	0010000011	–	1001000011	–	1000100011	–
47	1001000000011	100001000011	00000100011	00001000011	0010100011	0001000011	–	00100011	–	00010011
48	0001000000011	001000000011	01000010011	10010000011	01000000011	1000010011	–	00101011	–	0100000011
49	0010000000011	000001000011	01001010011	–	10000000011	–	–	0000000011	–	1000000011
50	0000010000011	010000100011	10001010011	–	10001000011	–	–	0000100011	–	1000100011
51	0100001000011	100000100011	10001001011	–	–	–	–	1001000011	–	–
52	1000001000011	100000010011	00000010011	00010000011	–	–	–	–	00100011	–
53	0000001000011	000000010011	00000000011	10000100011	–	0010000011	0001000011	–	001000011	–
54	0000000100011	000000010011	000010000011	–	00000000011	0010100011	1000010011	–	0000000011	–
55	0000100100011	000010010011	100100000011	–	00001000011	01000000011	–	–	0000100011	–
56	1001000100011	100100010011	–	00100000011	10010000011	10000000011	–	–	1001000011	–
57	0100000010011	000100100011	–	00000100011	–	10001000011	–	–	–	00100011
58	1000000010011	010000000011	000100000011	01000010011	–	–	–	–	–	001000011
59	0000000010011	1000000000011	100001000011	010000000011	–	–	–	0001000011	–	0000000011
60	0000100010011	1000100000011	–	100000000011	–	–	0010000011	1000010011	–	0000100011
61	1001000010011	0000000000011	001000000011	100010000011	00010000011	–	0010100011	–	–	1001000011
62	00010000010011	0000100000011	000001000011	–	10000100011	00000000011	01000000011	–	–	–
63	0010000010011	1001000000011	010000100011	00000010011	–	00001000011	10000000011	–	–	–
64	0000010010011	–	100000100011	000000000011	–	10010000011	10001000011	–	–	–
65	01000000000011	0001000000011	100000010011	000010000011	–	–	–	–	0001000011	–
66	10000000000011	1000010000011	100010010011	100100000011	00100000011	–	–	–	1000010011	–
67	00000000000011	0010000000011	000000100011	–	00000100011	–	–	0010000011	–	–
68	00001000000011	0000010000011	000000010011	–	01000010011	–	–	0010100011	–	–
69	10010000000011	0100001000011	000010010011	00010010011	010000000011	–	–	01000000011	–	–
70	00010000000011	1000001000011	100100010011	000100000011	100000000011	00010000011	00000000011	10000000011	–	–



**Table 4.** Continued

$k$	$(u = -11)$ $(v = -12)$	$(u = -12)$ $(v = -13)$	$(u = -13)$ $(v = -14)$	$(u = -14)$ $(v = -15)$	$(u = -15)$ $(v = -16)$	$(u = -16)$ $(v = -17)$	$(u = -17)$ $(v = -18)$	$(u = -18)$ $(v = -19)$	$(u = -19)$ $(v = -20)$	$(u = -20)$ $(v = 21)$
12	–	–	–	–	–	–	–	–	–	–
13	00011	–	–	–	–	–	–	–	–	–
14	1000011	00011	–	–	–	–	–	–	–	–
15	–	1000011	00011	–	–	–	–	–	–	–
16	–	–	1000011	00011	–	–	–	–	–	–
17	–	–	–	1000011	00011	–	–	–	–	–
18	–	–	–	–	1000011	00011	–	–	–	–
19	–	–	–	–	–	1000011	00011	–	–	–
20	–	–	–	–	–	–	1000011	00011	–	–
21	–	–	–	–	–	–	–	1000011	00011	–
22	–	–	–	–	–	–	–	–	1000011	00011
23	–	–	–	–	–	–	–	–	–	1000011
24	0011	–	–	–	–	–	–	–	–	–
25	0000011	–	–	–	–	–	–	–	–	–
26	01000011	0011	–	–	–	–	–	–	–	–
27	10000011	0000011	–	–	–	–	–	–	–	–
28	100000011	01000011	0011	–	–	–	–	–	–	–
29	100010011	10000011	0000011	–	–	–	–	–	–	–
30	–	100000011	01000011	0011	–	–	–	–	–	–
31	–	100010011	10000011	0000011	–	–	–	–	–	–
32	–	–	100000011	01000011	0011	–	–	–	–	–
33	–	–	100010011	10000011	0000011	–	–	–	–	–
34	–	–	–	100000011	01000011	0011	–	–	–	–
35	–	–	–	100010011	10000011	0000011	–	–	–	–
36	–	–	–	–	100000011	01000011	0011	–	–	–
37	–	–	–	–	100010011	10000011	0000011	–	–	–
38	00000011	–	–	–	–	100000011	01000011	0011	–	–
39	000000011	–	–	–	–	100010011	10000011	0000011	–	–
40	000010011	–	–	–	–	–	100000011	01000011	0011	–
41	100100011	00000011	–	–	–	–	100010011	10000011	0000011	–
42	–	000000011	–	–	–	–	–	100000011	01000011	0011
43	–	000010011	–	–	–	–	–	100010011	10000011	0000011
44	–	100100011	00000011	–	–	–	–	–	100000011	01000011
45	–	–	000000011	–	–	–	–	–	100010011	10000011
46	–	–	000010011	–	–	–	–	–	–	100000011
47	–	–	100100011	00000011	–	–	–	–	–	100010011
48	–	–	–	000000011	–	–	–	–	–	–
49	0010011	–	–	000010011	–	–	–	–	–	–
50	–	–	–	100100011	00000011	–	–	–	–	–
51	0100000011	–	–	–	000010011	–	–	–	–	–
52	1000000011	0010011	–	–	100100011	00000011	–	–	–	–
53	1000100011	–	–	–	–	000000011	–	–	–	–
54	–	00010011	–	–	–	000010011	–	–	–	–
55	–	0100000011	–	–	–	100100011	00000011	–	–	–
56	–	1000000011	0010011	–	–	–	000000011	–	–	–
57	–	1000100011	–	–	–	–	000010011	–	–	–
58	–	–	00010011	–	–	–	100100011	00000011	–	–

**Table 4.** Continued

$k$	$(u = -11)$ $(v = -12)$	$(u = -12)$ $(v = -13)$	$(u = -13)$ $(v = -14)$	$(u = -14)$ $(v = -15)$	$(u = 15)$ $(v = -16)$	$(u = -16)$ $(v = -17)$	$(u = -17)$ $(v = -18)$	$(u = -18)$ $(v = -19)$	$(u = -19)$ $(v = -20)$	$(u = -20)$ $(v = 21)$
59	-	-	010000011	-	-	-	-	000000011	-	-
60	-	-	100000011	0010011	-	-	-	000010011	-	-
61	00100011	-	1000100011	-	-	-	-	100100011	00000011	-
62	001000011	-	-	00010011	-	-	-	-	000000011	-
63	000000001 1	-	-	0100000011	-	-	-	-	000010011	-
64	000010001 1	-	-	1000000011	0010011	-	-	-	100100011	00000011
65	100100001 1	-	-	1000100011	-	-	-	-	-	000000011
66	-	00100011	-	-	00010011	-	-	-	-	000010011
67	-	001000011	-	-	0100000011	-	-	-	-	100100011
68	-	0000000011	-	-	1000000011	0010011	-	-	-	-
69	-	0000100011	-	-	1000100011	-	-	-	-	-
70	-	1001000011	-	-	-	00010011	-	-	-	-
71	-	-	00100011	-	-	0100000011	-	-	-	-
72	-	-	001000011	-	-	1000000011	0010011	-	-	-
73	-	-	0000000011	-	-	1000100011	-	-	-	-
74	-	-	0000100011	-	-	-	00010011	-	-	-
75	-	-	1001000011	-	-	-	0100000011	-	-	-
76	-	-	1001000011	-	-	-	0100000011	-	-	-
77	000100001 1	-	-	00100011	-	-	1000000011	0010011	-	-
78	-	-	-	001000011	-	-	1000100011	-	-	-
79	-	-	-	0000000011	-	-	-	00010011	-	-
80	-	-	-	0000100011	-	-	-	0100000011	-	-
81	-	-	-	1001000011	-	-	-	1000000011	0010011	-
82	-	-	-	-	00100011	-	-	1000100011	-	-
83	-	0001000011	-	-	001000011	-	-	-	00010011	-
84	-	-	-	-	0000000011	-	-	-	0100000011	-
85	-	-	-	-	0000100011	-	-	-	1000000011	0010011
86	-	-	-	-	1001000011	-	-	-	1000100011	-
87	-	-	-	-	-	00100011	-	-	-	00010011
88	001000001 1	-	-	-	-	001000011	-	-	-	0100000011
89	000001001 1	-	0001000011	-	-	0000000011	-	-	-	1000000011
90	010000000 11	-	-	-	-	0000100011	-	-	-	1000100011
91	100000000 11	-	-	-	-	1001000011	-	-	-	-
92	100010000 11	-	-	-	-	-	00100011	-	-	-
93	-	-	-	-	-	-	001000011	-	-	-
94	-	-	-	-	-	-	0000000011	-	-	-
95	-	0010000011	-	0001000011	-	-	0000100011	-	-	-
96	-	0000010011	-	-	-	-	1001000011	-	-	-

**Table 4.** Continued

$k$	$(u = -11)$ $(v = -12)$	$(u = -12)$ $(v = -13)$	$(u = -13)$ $(v = -14)$	$(u = -14)$ $(v = -15)$	$(u = -15)$ $(v = -16)$	$(u = -16)$ $(v = -17)$	$(u = -17)$ $(v = -18)$	$(u = -18)$ $(v = -19)$	$(u = -19)$ $(v = -20)$	$(u = -20)$ $(v = 21)$
97	–	0100000011	–	–	–	–	–	00100011	–	–
98	–	1000000011	–	–	–	–	–	00100011	–	–
99	–	1000100011	–	–	–	–	–	000000011	–	–
100	–	–	–	–	–	–	–	0000100011	–	–

Moreover, this section obtains the following theorem, which is the relation between the second-order bivariate variant of the Narayana sequence and the Narayana sequence.

**Theorem 3.2.3.** Let  $VN_{(u,v)}^{(2)}(k)$  is the second-order bivariate variant of the Narayana sequence and  $N(k)$  is the Narayana sequence. In this case, for  $k \geq 1$ ,

$$VN_{(u,v)}^{(2)}(k) = N(k - 3) - uN(k - 6) - v(N(k - 6) + N(k - 7))$$

where  $k$  is an integer.

PROOF. We have the following set from the second-order bivariate variant of the Narayana sequence.

$$VN_{(u,v)}^{(2)}(k) = \{u, v, 1 - u - v, 1 - v, 1, 2 - u - v, 3 - u - 2v, 4 - u - 2v, 6 - 2u - 3v, \dots\}$$

$$VN_{(u,v)}^{(2)}(1) = u = 0 - u(-1) - v((-1) + 1) = N(-2) - uN(-5) - v(N(-5) + N(-6))$$

$$VN_{(u,v)}^{(2)}(2) = v = 0 - u0 - v(0 - 1) = N(-1) - uN(-4) - v(N(-4) + N(-5))$$

The results are correct for  $k = 1, k = 2$ , and  $k = 3$ , as seen above. Suppose that the results are correct for  $k = 1, 2, \dots, n$ . In that case,

$$VN_{(u,v)}^{(2)}(n - 2) = N(n - 5) - uN(n - 8) - v(N(n - 8) + N(n - 9))$$

$$VN_{(u,v)}^{(2)}(n - 1) = N(n - 4) - uN(n - 7) - v(N(n - 7) + N(n - 8))$$

$$VN_{(u,v)}^{(2)}(n) = N(n - 3) - uN(n - 6) - v(N(n - 6) + N(n - 7))$$

Since

$$\begin{aligned} VN_{(u,v)}^{(2)}(n + 1) &= N(n - 5) - uN(n - 8) - v(N(n - 8) + N(n - 9)) + N(n - 4) - uN(n - 7) \\ &\quad - v(N(n - 7) + N(n - 8)) + N(n - 3) - uN(n - 6) - v(N(n - 6) + N(n - 7)) \\ &= N(n - 5) + N(n - 4) + N(n - 3) - u(N(n - 8) + N(n - 7) + N(n - 6)) \\ &\quad - v((N(n - 6) + N(n - 7) + N(n - 7)) + N(n - 8) + N(n - 8) + N(n - 9)) \\ &= N(n - 2) - uN(n - 5) - v(N(n - 5) + N(n - 6)) \\ &= N(n + 1 - 3) - uN(n + 1 - 6) - v(N(n + 1 - 6) + N(n + 1 - 7)). \end{aligned}$$

this equation is true for  $n + 1$ . Therefore, by the induction,  $VN_{(u,v)}^{(2)}(k) = N(k - 3) - uN(k - 6) - v(N(k - 6) + N(k - 7)), k \geq 1$ .



### 3.3. Cryptographic Comparison of Bivariate Gopala Hemachandra and Variant Narayana Codes

Cryptography is used to make sense of incomprehensible messages [14]. What cryptography focuses on is privacy. Cryptography aims to provide secure communication between two people. The information is named plaintext [15].

Cryptography is divided into two according to a key structure. These are symmetric cryptography and asymmetric cryptography. DES and AES are examples of symmetric cryptography, while RSA and ElGamal are examples of asymmetric cryptography. Different cryptographic systems are also built, too. One of them is the system obtained using source coding. A few examples made with the system can be found in [10], [12]. Hence, new source codes of third and fourth order identified in this study can also be used in cryptographic applications. But columns with undetectable codes can't be used in cryptography. While for the third-order bivariate GH code, there are no undetectable codes in  $GH_{(-1,-2)}^{(3)}(k)$ ,  $GH_{(-2,-3)}^{(3)}(k)$ ,  $GH_{(-3,-4)}^{(3)}(k)$  columns, for the second-order bivariate variant of the Narayana code, there are no undetectable codes in only  $VN_{(-1,-2)}^{(2)}(k)$  column. In addition, for the positive integers  $k = 1, 2, 3$ , the third-order bivariate GH code  $GH_{(u,v)}^{(3)}(k)$  precisely exists, for a bivariate variant of Narayana code, for the positive integers  $k = 1$ , the second order bivariate variant of Narayana code  $VN_{(u,v)}^{(2)}(k)$  exactly exists. Similarly, for  $1 \leq k \leq 100$ , while there is at most  $j$  consecutive undetectable (--) values, the third-order bivariate GH code in  $GH_{(-(3+j),-(4+j))}^{(3)}(k)$  column in which  $1 \leq j \leq 17$ , there is at most  $j$  consecutive undetectable (--) values the second order bivariate variant of Narayana code in  $VN_{(-(1+j),-(2+j))}^{(2)}(k)$  column in which  $1 \leq j \leq 19$ . At that rate, we obtained that the third-order bivariate variant of the Fibonacci code is more valuable than the second-order bivariate variant of the Narayana code in terms of cryptography.

## 4. Conclusion

In this study, firstly, we examined Fibonacci, Tribonacci, and Narayana sequences and investigated the third-order variant of the Fibonacci sequence, and also obtained which is the relation between the third-order variant of the Fibonacci sequence  $VF_u^{(3)}(k)$  and the Tribonacci sequence  $T(k)$  for any integer  $k \geq 1$  with a theorem. Then, we described a new variant of the Fibonacci sequence and a new variant of the Narayana sequence, depending on two negative integer variables  $u$  and  $v$ . Moreover, we named these new sequences the third-order bivariate variant of the Fibonacci sequence and the second-order bivariate variant of the Narayana sequence, respectively. Then, we obtained a bivariate variant of the GH code and a bivariate variant of the Narayana universal code based on these bivariate variant sequences we described.

In addition, we obtained the relation between the third-order bivariate variant of the Fibonacci sequence  $VF_{(u,v)}^{(3)}(k)$  and the Tribonacci sequence  $T(k)$ , and the second-order bivariate variant of the Narayana sequence  $VN_{(u,v)}^{(2)}(k)$  and the Narayana sequence  $N(k)$  for any integer  $k \geq 1$ .

Afterwards, we showed in tables  $VF_{(u,v)}^{(3)}(k)$ ,  $VN_{(u,v)}^{(2)}(k)$  we have defined for  $1 \leq k \leq 100$  and  $u = -1, -2, \dots, -20$  and  $v = -2, -3, \dots, -21$  ( $u$  and  $v$  are consecutive,  $v < u$ ). We got some important results from the tables for these bivariate variants of universal codes. For  $k = 1, 2, 3$ , the third-order bivariate GH code  $GH_{(u,v)}^{(3)}(k)$  exactly exists. There are at most  $j$  consecutive undetectable (-) values in the third-order bivariate GH code in  $GH_{(-(3+j),-(4+j))}^{(3)}(k)$  column in which  $1 \leq j \leq 17$  and for  $1 \leq k \leq 100$ . As long as  $j$  raises, the detectable of GH code is reduced in  $GH_{(-(3+j),-(4+j))}^{(3)}(k)$  column in which  $1 \leq j \leq 17$  and for  $1 \leq k \leq 100$ . For  $k = 1$ , the second-order bivariate variant of the Narayana code  $VN_{(u,v)}^{(2)}(k)$  exactly exists. For

$1 \leq k \leq 100$ , there are at most  $j$  consecutive undetectable (--) values in the second-order bivariate variant of Narayana code in  $VN_{(-(1+j), -(2+j))}^{(2)}(k)$  column in which  $1 \leq j \leq 19$ . And for  $1 \leq k \leq 100$ , as long as  $j$  raises, the detectable of Narayana code is reduced in  $VN_{(-(1+j), -(2+j))}^{(2)}(k)$  column in which  $1 \leq j \leq 19$ .

Finally, we compared the two universal codes defined here regarding the existence of these codes for positive integers  $k$  and the use of these sequences in cryptography. And, according to the above results, in cryptographic applications, columns  $GH_{(-1, -2)}^{(3)}(k)$ ,  $GH_{(-2, -3)}^{(3)}(k)$ ,  $GH_{(-3, -4)}^{(3)}(k)$  can be precisely used for a bivariate GH code for  $m = 3$ , the column  $VN_{(-1, -2)}^{(2)}(k)$  can be precisely used for a bivariate variant of the Narayana code for  $m = 2$ . And hence, we acquired the third-order bivariate variant of the Fibonacci code is more valuable than the second-order bivariate variant of the Narayana code.

## Author Contributions

The author read and approved the last version of the manuscript.

## Conflict of Interest

The author declares no conflict of interest.

## References

- [1] T. Koshy, *Fibonacci and Lucas Numbers with Applications*, Wiley, New York, 2001.
- [2] M. Feinberg, *Fibonacci-Tribonacci*, *The Fibonacci Quarterly* 1 (1963) 71–74.
- [3] J. H. Thomas, *Variations on the Fibonacci Universal Code*. arXiv:0701085 (2007).
- [4] S. Kak, *The Golden Mean and the Physics of Aesthetics*, in: B. Yadav, M. Mohan (Eds.), *Ancient Indian Leaps into Mathematics*, Birkhäuser, Boston, 2011, pp. 111–119.
- [5] K. Kirthi, S. Kak, *The Narayana Universal Code*, arXiv: 1601.07110 (2016).
- [6] T. Buschmann, L.V. Bystrykh, *Levenshtein Error-Correcting Barcodes for Multiplexed DNA Sequencing*, *BMC Bioinformatics* 14 (1) (2013) 272.
- [7] E. Zeckendorf, *Representation Des Nombres Naturels Par Une Somme Des Nombres De Fibonacci Ou De Nombres De Lucas*, *Bulletin De La Society Royale des Sciences de Liege* 41 (1972) 179–182.
- [8] S. T. Klein, M. K. Ben-Nissan, *On the Usefulness of Fibonacci Compression Codes*, *The Computer Journal* 53 (6) (2010) 701–716.
- [9] M. Basu, B. Prasad, *Long Range Variant of Fibonacci Universal Code*, *Journal of Number Theory* 130 (2010) 1925–1931.
- [10] A. Nalli, Ç. Özyılmaz, *The Third order Variations on the Fibonacci Universal Code*, *Journal of Number Theory* 149 (2015) 15–32.
- [11] D.E. Daykin, *Representation of Natural Numbers as Sums of Generalized Fibonacci Numbers*, *Journal of London Mathematical Society* 35 (1960) 143–160.
- [12] M. Basu, M. Das, *Uses of Second Order Variant Fibonacci Universal Code in Cryptography*, *Control and Cybernetics* 45 (2) (2016) 239–257.
- [13] M. Das, S. Sinha, *A Variant of the Narayana Coding Scheme*, *Control and Cybernetics* 48 (3) (2019) 473–484.
- [14] C. Çimen, S. Akleyek, E. Akyıldız, *Şifrelerin Matematiği Kriptografi*, METU Press Ankara, 2007.
- [15] D. R. Stinson, *Cryptography Theory and Practice*, Chapman & Hall, Ohio, CRC Press, 2002.